Chapter 1

Generalized Information Measures for Arbitrary Systems with Memory

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Shannon's entropy

• Entropy of a discrete random variable X:

$$H(X) := -\sum_{x \in \mathcal{X}} P_X(x) \log P_X(x) = E_X \left[-\log P_X(X) \right] \text{ nats}$$

is a measure of the average amount of uncertainty in X.

• Entropy rate for a sequence of random variables $X_1, X_2, \ldots, X_n, \ldots$ is

$$\lim_{n \to \infty} \frac{1}{n} H(X^n) = \lim_{n \to \infty} \frac{1}{n} E\left[-\log P_{X^n}(X^n)\right],$$

assuming the limit exists.

- Operation meaning: Shannon's source coding theorems for stationary and ergodic systems.
- Question: Does these measures have the same operational significance for systems with time-varying and non-stationary statistics. Answer: No.
- Solution: Require new entropy measure which can appropriately characterize the operational limits of *arbitrary* stochastic systems.

- In general, there are two indices for random variables or observations: a time index and a space index.
- When a sequence of random variables is denoted by

$$X_1, X_2, \ldots, X_n, \ldots,$$

the subscript i of X_i can be treated as either a time index or a space index, but not both.

• Hence, when a sequence of random variables is a function of both time and space, the notation of $X_1, X_2, \ldots, X_n, \ldots$, is by no means sufficient; and therefore, a new model for a general time-varying source, such as

$$X_1^{(n)}, X_2^{(n)}, \dots, X_t^{(n)}, \dots,$$

where t is the time index and n is the space or position index (or vice versa), becomes significant.

• When block-wise (fixed-length) compression of such source (with blocklength n) is considered, the same question as to the compression of i.i.d. source arises:

what is the minimum compression rate (say in bits per source sample) for which the probability of error probability can be made arbitrarily small as the blocklength goes to infinity?

• To answer the question, information theorists have to find a sequence of data compression codes for each blocklength *n* and investigate if the decompression error goes to zero as *n* approaches infinity.

• However, unlike those simple source models such as discrete memorylessness, the source being arbitrary may exhibit distinct statistics for each blocklength n; e.g., for

$$n = 1 : X_{1}^{(1)}$$

$$n = 2 : X_{1}^{(2)}, X_{2}^{(2)}$$

$$n = 3 : X_{1}^{(3)}, X_{2}^{(3)}, X_{3}^{(3)}$$

$$n = 4 : X_{1}^{(4)}, X_{2}^{(4)}, X_{3}^{(4)}, X_{4}^{(4)}$$
(1.0.2)

the statistics of $X_1^{(4)}$ could be different from $X_1^{(1)}$, $X_1^{(2)}$ and $X_1^{(3)}$ (i.e., the source statistics are not necessarily consistent).

• Since the model in question (1.0.2) is general, and the system statistics can be *arbitrarily* defined, it is therefore named an *arbitrary system with memory*.

• The triangular array of random variables in (1.0.2) is denoted by a boldface letter as

$$\boldsymbol{X} := \{X^n\}_{n=1}^{\infty},$$

where

$$X^{n} := \left(X_{1}^{(n)}, X_{2}^{(n)}, \dots, X_{n}^{(n)} \right);$$

for convenience, we also write

$$\boldsymbol{X} := \left\{ X^n = \left(X_1^{(n)}, X_2^{(n)}, \dots, X_n^{(n)} \right) \right\}_{n=1}^{\infty}.$$

Spectrum and Quantile

Definition 1.1 (Inf/sup-spectrum) If $\{A_n\}_{n=1}^{\infty}$ is a sequence of random variables, then its *inf-spectrum* $\underline{u}(\cdot)$ and its *sup-spectrum* $\overline{u}(\cdot)$ are defined by

$$\underline{u}(\theta) := \liminf_{n \to \infty} \Pr\{A_n \le \theta\} \text{ and } \overline{u}(\theta) := \limsup_{n \to \infty} \Pr\{A_n \le \theta\},\$$

respectively, where $\theta \in \mathbb{R}$.

• $\underline{u}(\cdot)$ and $\overline{u}(\cdot)$ are respectively the limit and the limit of the cumulative distribution function (CDF) of A_n .

Definition 1.2 (Quantile of inf/sup-spectrum) For any $0 \leq \delta \leq 1$, For any $0 \leq \delta \leq 1$, the quantile \underline{U}_{δ} of the sup-spectrum $\overline{u}(\cdot)$ and the quantile \overline{U}_{δ} of the inf-spectrum $\underline{u}(\cdot)$ are defined by

$$\underline{U}_{\delta} := \sup\{\theta : \overline{u}(\theta) \le \delta\} \text{ and } \overline{U}_{\delta} := \sup\{\theta : \underline{u}(\theta) \le \delta\},\$$

respectively. It follows from the above definitions that \underline{U}_{δ} and \overline{U}_{δ} are right-continuous and non-decreasing in δ . Note that the supremum of an empty set is defined to be $-\infty$.

• If $\bar{u}(\cdot)$ is strictly increasing, then the quantile is exactly its inverse: $\underline{U}_{\delta} = \bar{u}^{-1}(\delta)$.

Liminf in probability and limsup in probability II: 1-7

• limit in probability \underline{U} of $\{A_n\}_{n=1}^{\infty}$ is the largest extended real number such that for all $\xi > 0$,

$$\lim_{n \to \infty} \Pr[A_n \le \underline{U} - \xi] = 0.$$

• limsup in probability \overline{U} of $\{A_n\}_{n=1}^{\infty}$ is the smallest extended real number such that for all $\xi > 0$,

 $\lim_{n \to \infty} \Pr[A_n \ge \bar{U} + \xi] = 0.$

$$\underline{U} = \lim_{\delta \downarrow 0} \underline{U}_{\delta} = \underline{U}_{0}$$

and

$$\bar{U} = \lim_{\delta \uparrow 1} \bar{U}_{\delta} = \sup\{\theta : \underline{u}(\theta) < 1\}.$$

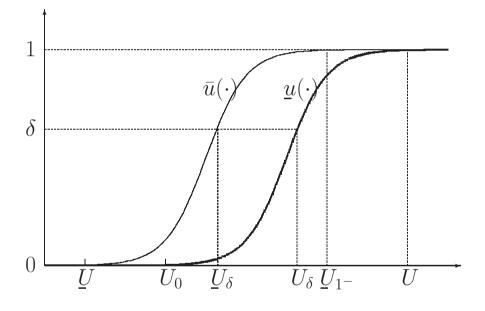
• It readily follows from the above definitions that

$$\underline{U} \leq \underline{U}_{\delta} \leq \overline{U}_{\delta} \leq \overline{U} \quad \text{for } \delta \in [0, 1).$$

• $\bar{U}_1 = \underline{U}_1 = \infty.$

Liminf in probability and limsup in probability

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The asymptotic CDFs (spectrums) of $\{A_n\}_{n=1}^{\infty}$ and their quantiles: $\bar{u}(\cdot) = \text{sup-spectrum of } \{A_n\}, \quad \underline{u}(\cdot) = \text{inf-spectrum of } \{A_n\},$ $\underline{U}_{\delta} = \text{quantile of } \bar{u}(\cdot), \quad \overline{U}_{\delta} = \text{quantile of } \underline{u}(\cdot),$ $\underline{U} = \lim_{\delta \downarrow 0} \underline{U}_{\delta} = \underline{U}_0, \quad \underline{U}_{1^-} = \lim_{\xi \uparrow 1} \underline{U}_{\xi}, \quad \overline{U} = \lim_{\delta \uparrow 1} \overline{U}_{\delta}.$

Properties of quantile

Lemma 1.4 Assume:

- Two random sequences: $\{A_n\}_{n=1}^{\infty}$ and $\{B_n\}_{n=1}^{\infty}$;
- $\bar{u}(\cdot) = \text{sup-spectrum of } \{A_n\}_{n=1}^{\infty}; \underline{U}_{\delta} = \text{quantile of } \bar{u}(\cdot);$
- $\underline{u}(\cdot) = \text{inf-spectrum of } \{A_n\}_{n=1}^{\infty}; \, \overline{U}_{\delta} = \text{quantile of } \underline{u}(\cdot);$
- $\bar{v}(\cdot) = \text{sup-spectrum of } \{B_n\}_{n=1}^{\infty}; \underline{V}_{\delta} = \text{quantile of } \bar{u}(\cdot);$
- $\underline{v}(\cdot) = \text{inf-spectrum of } \{B_n\}_{n=1}^{\infty}; \, \overline{V}_{\delta} = \text{quantile of } \underline{u}(\cdot);$

•
$$(\overline{u+v})(\cdot) =$$
sup-spectrum of $\{A_n + B_n\}_{n=1}^{\infty}$, i.e.,

$$(\overline{u+v})(\theta) := \limsup_{n \to \infty} \Pr\{A_n + B_n \le \theta\};$$

 $(\underline{U+V})_{\delta} =$ quantile of $(\overline{u+v})(\cdot);$

•
$$(\underline{u}+\underline{v})(\cdot) = \text{inf-spectrum of } \{A_n + B_n\}_{n=1}^{\infty}, \text{ i.e.,}$$

 $(\underline{u}+\underline{v})(\theta) := \liminf_{n \to \infty} \Pr\{A_n + B_n \le \theta\};$
 $(\overline{U}+\overline{V})_{\delta} = \text{quantile of } (\underline{u}+\underline{v})(\cdot).$

Properties of quantile

Then the following statements hold.

- 1. \underline{U}_{δ} and \overline{U}_{δ} are both non-decreasing and right-continuous functions of δ for $\delta \in [0, 1]$.
- 2. $\lim_{\delta \downarrow 0} \underline{U}_{\delta} = \underline{U}_0$ and $\lim_{\delta \downarrow 0} \overline{U}_{\delta} = \overline{U}_0$.
- 3. For $\delta \ge 0$, $\gamma \ge 0$, and $\delta + \gamma \le 1$,

$$(\underline{U+V})_{\delta+\gamma} \ge \underline{U}_{\delta} + \underline{V}_{\gamma}, \qquad (1.2.1)$$

and

$$(\overline{U+V})_{\delta+\gamma} \ge \underline{U}_{\delta} + \overline{V}_{\gamma}.$$
(1.2.2)

4. For $\delta \ge 0, \gamma \ge 0$, and $\delta + \gamma \le 1$,

$$(\underline{U+V})_{\delta} \le \underline{U}_{\delta+\gamma} + \overline{V}_{(1-\gamma)}, \qquad (1.2.3)$$

and

$$(\overline{U+V})_{\delta} \le \bar{U}_{\delta+\gamma} + \bar{V}_{(1-\gamma)}.$$
(1.2.4)

<u>Generalized information measures</u>

In Definitions 1.1 and 1.2,

• $A_n = normalized entropy density$, i.e.,

$$\frac{1}{n}h_{X^n}(X^n) := -\frac{1}{n}\log P_{X^n}(X^n),$$

 δ -inf-entropy rate $\underline{H}_{\delta}(\mathbf{X})$ = quantile of the sup-spectrum of $\frac{1}{n}h_{X^n}(X^n)$

 δ -sup-entropy rate $\bar{H}_{\delta}(\mathbf{X})$ = quantile of the inf-spectrum of $\frac{1}{n}h_{X^n}(X^n)$.

•
$$A_n = normalized information density$$
, i.e.,

$$\frac{1}{n}i_{X^nW^n}(X^n;Y^n) = \frac{1}{n}i_{X^n,Y^n}(X^n;Y^n) := \frac{1}{n}\log\frac{P_{X^n,Y^n}(X^n,Y^n)}{P_{X^n}(X^n)P_{Y^n}(Y^n)},$$

 δ -inf-information-rate $I_{\delta}(\mathbf{X}; \mathbf{Y})$ = quantile of the sup-spectrum of $\frac{1}{n} i_{X^n W^n}(X^n; Y^n)$

 δ -sup-information-rate $\bar{I}_{\delta}(\boldsymbol{X};\boldsymbol{Y}) =$ quantile of the inf-spectrum of $\frac{1}{n}i_{X^{n}W^{n}}(X^{n};Y^{n}).$

<u>Generalized information measures</u>

• $A_n = normalized \ log-likelihood \ ratio, i.e.,$

$$\frac{1}{n} d_{X^n}(X^n \| \hat{X}^n) := \frac{1}{n} \log \frac{P_{X^n}(X^n)}{P_{\hat{X}^n}(X^n)}$$

 δ -inf-divergence rate $\underline{D}_{\delta}(\boldsymbol{X} \| \hat{\boldsymbol{X}})$ = quantile of the sup-spectrum of $\frac{1}{n} d_{X^n}(X^n \| \hat{X}^n)$

 δ -sup-divergence rate $\bar{D}_{\delta}(\boldsymbol{X} \| \hat{\boldsymbol{X}})$ = quantile of the inf-spectrum of $\frac{1}{n} d_{X^n}(X^n \| \hat{X}^n)$.

- Notes:
 - The *inf-entropy-rate* $\underline{H}(\mathbf{X})$ and the *sup-entropy-rate* $\overline{H}(\mathbf{X})$ are special cases of the δ -inf/sup-entropy rate measures:

$$\underline{H}(\mathbf{X}) = \underline{H}_0(\mathbf{X}), \text{ and } \overline{H}(\mathbf{X}) = \lim_{\delta \uparrow 1} \overline{H}_\delta(\mathbf{X}).$$

- Concept: If the random variable $(1/n)h(X^n)$ exhibits a limiting distribution, and suppose the limiting distribution of $(1/n)h_{X^n}(X^n)$ is positive over (-2, 2); and zero, otherwise. Then $\bar{H}(\mathbf{X}) = 2$ and $\underline{H}(\mathbf{X}) = -2$.

Generalized information measures

Entropy Measures		
system	arbitrary source \boldsymbol{X}	
norm. entropy density	$\frac{1}{n}h_{X^n}(X^n) := -\frac{1}{n}\log P_{X^n}(X^n)$	
entropy sup-spectrum	$\left \bar{h}(\theta) := \limsup_{n \to \infty} \Pr\left\{ \frac{1}{n} h_{X^n}(X^n) \le \theta \right\} \right $	
entropy inf-spectrum	$\underline{h}(\theta) := \liminf_{n \to \infty} \Pr\left\{\frac{1}{n} h_{X^n}(X^n) \le \theta\right\}$	
δ -inf-entropy rate	$\underline{H}_{\delta}(\boldsymbol{X}) := \sup\{\theta : \overline{h}(\theta) \le \delta\}$	
δ -sup-entropy rate	$\bar{H}_{\delta}(\boldsymbol{X}) := \sup\{\boldsymbol{\theta} : \underline{h}(\boldsymbol{\theta}) \le \delta\}$	
sup-entropy rate	$\bar{H}(\boldsymbol{X}) := \lim_{\delta \uparrow 1} \bar{H}_{\delta}(\boldsymbol{X})$	
inf-entropy rate	$\underline{H}(\boldsymbol{X}):=\underline{H}_0(\boldsymbol{X})$	

Generalized entropy measures where $\delta \in [0, 1]$.

Generalized information measures

Mutual Information Measures	
system	arbitrary channel $P_{W} = P_{Y X}$
	with input \boldsymbol{X} and output \boldsymbol{Y}
norm. information density	$\frac{1}{n}i_{X^nW^n}(X^n;Y^n)$
	$:= \frac{1}{n} \log \frac{P_{X^{n}, Y^{n}}(X^{n}, Y^{n})}{P_{X^{n}}(X^{n}) \times P_{Y^{n}}(Y^{n})}$
information sup-spectrum	$\left \overline{i}(\theta) := \limsup_{n \to \infty} \Pr\left\{ \frac{1}{n} i_{X^n W^n}(X^n; Y^n) \le \theta \right\} \right $
information inf-Spectrum	$\underline{i}(\theta) := \liminf_{n \to \infty} \Pr\left\{\frac{1}{n} i_{X^n W^n}(X^n; Y^n) \le \theta\right\}$
δ -inf-information rate	$I_{\delta}(\boldsymbol{X};\boldsymbol{Y}) := \sup\{\boldsymbol{\theta} : \bar{i}(\boldsymbol{\theta}) \le \delta\}$
δ -Sup-Information Rate	$\bar{I}_{\delta}(\boldsymbol{X};\boldsymbol{Y}) := \sup\{\theta : \underline{i}(\theta) \le \delta\}$
sup-information rate	$\bar{I}(\boldsymbol{X};\boldsymbol{Y}) := \lim_{\delta \uparrow 1} \bar{I}_{\delta}(\boldsymbol{X};\boldsymbol{Y})$
inf-information rate	$\underline{I}(\boldsymbol{X};\boldsymbol{Y}) := \underline{I}_0(\boldsymbol{X};\boldsymbol{Y})$

Generalized mutual information measures where $\delta \in [0, 1]$.

<u>Generalized information measures</u>

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Divergence Measures	
system	arbitrary sources \boldsymbol{X} and $\hat{\boldsymbol{X}}$
norm. log-likelihood ratio	$\frac{1}{n} d_{X^n}(X^n \ \hat{X}^n) := \frac{1}{n} \log \frac{P_{X^n}(X^n)}{P_{\hat{X}^n}(X^n)}$
divergence sup-spectrum	$\left \bar{d}(\theta) := \limsup_{n \to \infty} \Pr\left\{ \frac{1}{n} d_{X^n}(X^n \ \hat{X}^n) \le \theta \right\} \right $
divergence inf-spectrum	$\underline{d}(\theta) := \liminf_{n \to \infty} \Pr\left\{\frac{1}{n} d_{X^n}(X^n \ \hat{X}^n) \le \theta\right\}$
δ -inf-divergence rate	$\underline{D}_{\delta}(\boldsymbol{X} \ \hat{\boldsymbol{X}}) := \sup \{ \theta : \overline{d}(\theta) \le \delta \}$
δ -sup-divergence rate	$\bar{D}_{\delta}(\boldsymbol{X} \ \hat{\boldsymbol{X}}) := \sup\{\theta : \underline{d}(\theta) \leq \delta\}$
sup-divergence rate	$\bar{D}(\boldsymbol{X} \ \hat{\boldsymbol{X}}) := \lim_{\delta \uparrow 1} \bar{D}_{\delta}(\boldsymbol{X} \ \hat{\boldsymbol{X}})$
inf-divergence rate	$\underline{D}(\boldsymbol{X} \ \hat{\boldsymbol{X}}) := \underline{D}_0(\boldsymbol{X} \ \hat{\boldsymbol{X}})$

Generalized divergence measures where $\delta \in [0, 1]$.

• An example of basic properties for Shannon's entropy: I(X;Y) = H(Y) - H(Y|X).

– By taking $\delta = 0$ and letting $\gamma \downarrow 0$ in

 $(\underline{U+V})_{\delta+\gamma} \ge \underline{U}_{\delta} + \underline{V}_{\gamma}$ for $\delta \ge 0, \gamma \ge 0$, and $\delta + \gamma \le 1$

and

$$(\underline{U+V})_{\delta} \leq \underline{U}_{\delta+\gamma} + \overline{V}_{(1-\gamma)} \text{ for } \delta \geq 0, \gamma \geq 0, \text{ and } \delta+\gamma \leq 1,$$

we obtain

$$(\underline{U+V}) \ge \underline{U}_0 + \lim_{\gamma \downarrow 0} \underline{V}_\gamma \ge \underline{U} + \underline{V}$$

and

$$(\underline{U+V}) \le \lim_{\gamma \downarrow 0} \underline{U}_{\gamma} + \lim_{\gamma \downarrow 0} \overline{V}_{(1-\gamma)} = \underline{U} + \overline{V}.$$

- Meaning: The limit in probability of a sequence of random variables $A_n + B_n$ is upper bounded by the limit in probability of A_n plus the limit in probability of B_n ; and is lower bounded by the sum of the limit in probability of A_n and B_n .

– This fact can be used to show that

$$\underline{I}(\boldsymbol{X};\boldsymbol{Y}) + \underline{H}(\boldsymbol{Y}|\boldsymbol{X}) \leq \underline{H}(\boldsymbol{Y}) \leq \underline{I}(\boldsymbol{X};\boldsymbol{Y}) + \overline{H}(\boldsymbol{Y}|\boldsymbol{X}),$$

or equivalently,

$$\underline{H}(\boldsymbol{Y}) - \overline{H}(\boldsymbol{Y}|\boldsymbol{X}) \leq \underline{I}(\boldsymbol{X};\boldsymbol{Y}) \leq \underline{H}(\boldsymbol{Y}) - \underline{H}(\boldsymbol{Y}|\boldsymbol{X}).$$

Lemma 1.5 For a source X with finite alphabet \mathcal{X} and arbitrary sources Y and Z, the following properties hold.

- 1. $\bar{H}_{\delta}(\boldsymbol{X}) \geq 0$ for $\delta \in [0, 1]$. (This property also applies to $\underline{H}_{\delta}(\boldsymbol{X})$, $\bar{I}_{\delta}(\boldsymbol{X}; \boldsymbol{Y})$, $\underline{I}_{\delta}(\boldsymbol{X}; \boldsymbol{Y}), \ \bar{D}_{\delta}(\boldsymbol{X} \| \hat{\boldsymbol{X}})$, and $\underline{D}_{\delta}(\boldsymbol{X} \| \hat{\boldsymbol{X}})$.)
- 2. $I_{\delta}(\boldsymbol{X};\boldsymbol{Y}) = I_{\delta}(\boldsymbol{Y};\boldsymbol{X})$ and $\bar{I}_{\delta}(\boldsymbol{X};\boldsymbol{Y}) = \bar{I}_{\delta}(\boldsymbol{Y};\boldsymbol{X})$ for $\delta \in [0,1]$.

3. For $0 \leq \delta < 1, 0 \leq \gamma < 1$ and $\delta + \gamma \leq 1$,

$$\underline{I}_{\delta}(\boldsymbol{X};\boldsymbol{Y}) \leq \underline{H}_{\delta+\gamma}(\boldsymbol{Y}) - \underline{H}_{\gamma}(\boldsymbol{Y}|\boldsymbol{X}), \qquad (1.4.1)$$

$$\underline{I}_{\delta}(\boldsymbol{X};\boldsymbol{Y}) \leq \bar{H}_{\delta+\gamma}(\boldsymbol{Y}) - \bar{H}_{\gamma}(\boldsymbol{Y}|\boldsymbol{X}), \qquad (1.4.2)$$

$$\bar{I}_{\delta}(\boldsymbol{X};\boldsymbol{Y}) \leq \bar{H}_{\delta+\gamma}(\boldsymbol{Y}) - \underline{H}_{\gamma}(\boldsymbol{Y}|\boldsymbol{X}), \qquad (1.4.3)$$

$$\underline{I}_{\delta+\gamma}(\boldsymbol{X};\boldsymbol{Y}) \geq \underline{H}_{\delta}(\boldsymbol{Y}) - \bar{H}_{(1-\gamma)}(\boldsymbol{Y}|\boldsymbol{X}), \qquad (1.4.4)$$

and

$$\bar{I}_{\delta+\gamma}(\boldsymbol{X};\boldsymbol{Y}) \geq \bar{H}_{\delta}(\boldsymbol{Y}) - \bar{H}_{(1-\gamma)}(\boldsymbol{Y}|\boldsymbol{X}).$$
(1.4.5)

(Note that the case of $(\delta, \gamma) = (1, 0)$ holds for (1.4.1) and (1.4.2), and the case of $(\delta, \gamma) = (0, 1)$ holds for (1.4.3), (1.4.4) and (1.4.5).)

4. $0 \leq \underline{H}_{\delta}(\mathbf{X}) \leq \overline{H}_{\delta}(\mathbf{X}) \leq \log |\mathcal{X}|$ for $\delta \in [0, 1)$, where each $X_i^{(n)}$ takes values in \mathcal{X} for $i = 1, \ldots, n$ and $n = 1, 2, \ldots$.

5. $I_{\delta}(\boldsymbol{X}, \boldsymbol{Y}; \boldsymbol{Z}) \geq I_{\delta}(\boldsymbol{X}; \boldsymbol{Z})$ for $\delta \in [0, 1]$.

Property 1:

$$\Pr\left\{-\frac{1}{n}\log P_{X^{n}}(X^{n}) < 0\right\} = 0,$$

$$\Pr\left\{\frac{1}{n}\log\frac{P_{X^{n}}(X^{n})}{P_{\hat{X}^{n}}(X^{n})} < -\nu\right\} = P_{X^{n}}\left\{x^{n} \in \mathcal{X}^{n} : \frac{1}{n}\log\frac{P_{X^{n}}(x^{n})}{P_{\hat{X}^{n}}(x^{n})} < -\nu\right\}$$

$$= \sum_{\left\{x^{n} \in \mathcal{X}^{n} : P_{X^{n}}(x^{n}) < P_{\hat{X}^{n}}(x^{n})e^{-n\nu}\right\}}$$

$$\leq \sum_{\left\{x^{n} \in \mathcal{X}^{n} : P_{X^{n}}(x^{n}) < P_{\hat{X}^{n}}(x^{n})e^{n\nu}\right\}}$$

$$\leq e^{-n\nu} \cdot \sum_{\left\{x^{n} \in \mathcal{X}^{n} : P_{X^{n}}(x^{n}) < P_{\hat{X}^{n}}(x^{n})e^{n\nu}\right\}}$$

$$\leq e^{-\nu n}, \qquad (1.4.6)$$

and, by following the same procedure as (1.4.6),

$$\Pr\left\{\frac{1}{n}\log\frac{P_{X^{n},Y^{n}}(X^{n},Y^{n})}{P_{X^{n}}(X^{n})P_{Y^{n}}(Y^{n})} < -\nu\right\} \leq e^{-\nu n}.$$

Property 2: An immediate consequence of the definition.

Property 3: Follow from the facts that

$$\frac{1}{n}h_{Y^n}(Y^n) = \frac{1}{n}i_{X^n,Y^n}(X^n;Y^n) + \frac{1}{n}h_{X^n,Y^n}(Y^n|X^n),$$

where

$$\frac{1}{n}h_{X^n,Y^n}(Y^n|X^n) := -\frac{1}{n}\log P_{Y^n|X^n}(Y^n|X^n).$$

Property 4: $\bar{H}_{\delta}(\cdot)$ is non-decreasing in δ , $\bar{H}_{\delta}(\mathbf{X}) \leq \bar{H}(\mathbf{X})$, and $\bar{H}(\mathbf{X}) \leq \log |\mathcal{X}|$. The last inequality can be proved as follows:

$$\Pr\left\{\frac{1}{n}h_{X^{n}}(X^{n}) \leq \log|\mathcal{X}| + \nu\right\} |$$

= $1 - P_{X^{n}}\left\{x^{n} \in \mathcal{X}^{n} : \frac{1}{n}\log\frac{P_{X^{n}}(X^{n})}{1/|\mathcal{X}|^{n}} < -\nu\right\}$
 $\geq 1 - e^{-n\nu},$

where the last step can be obtained by using the same procedure as (1.4.6). Therefore, $\underline{h}(\log |\mathcal{X}| + \nu) = 1$ for any $\nu > 0$, which indicates that $\overline{H}(\mathbf{X}) \leq \log |\mathcal{X}|$.

Property 5:

$$\frac{1}{n}i_{X^n,Y^n,Z^n}(X^n,Y^n;Z^n) = \frac{1}{n}i_{X^n,Z^n}(X^n;Z^n) + \frac{1}{n}i_{X^n,Y^n,Z^n}(Y^n;Z^n|X^n).$$

Lemma 1.6 (Data processing lemma) Fix $\delta \in [0, 1]$. Consider a channel with finite input and output alphabets, \mathcal{X} and \mathcal{Y} , respectively, and with distribution $P_{W^n}(y^n|x^n) = P_{Y^n|X^n}(y^n|x^n) = \prod_{i=1}^n P_{Y_i|X_i}(y_i|x_i)$ for all $n, x^n \in \mathcal{X}^n$ and $y^n \in \mathcal{Y}^n$. Then

$$\underline{I}_{\delta}(\boldsymbol{X}_1; \boldsymbol{X}_3) \leq \underline{I}_{\delta}(\boldsymbol{X}_1; \boldsymbol{X}_2).$$

Proof: By Property 5 of Lemma 1.5, we get

$$\underline{I}_{\delta}(\boldsymbol{X}_1; \boldsymbol{X}_3) \leq \underline{I}_{\delta}(\boldsymbol{X}_1; \boldsymbol{X}_2, \boldsymbol{X}_3) = \underline{I}_{\delta}(\boldsymbol{X}_1; \boldsymbol{X}_2),$$

where the equality holds because

$$\frac{1}{n}\log\frac{P_{X_1^n,X_2^n,X_3^n}(x_1^n,x_2^n,x_3^n)}{P_{X_1^n}(x_1^n)P_{X_2^n,X_3^n}(x_2^n,x_3^n)} = \frac{1}{n}\log\frac{P_{X_1^n,X_2^n}(x_1^n,x_2^n)}{P_{X_1^n}(x_1^n)P_{X_2^n}(x_2^n)}.$$

Lemma 1.7 (Optimality of independent inputs) Fix $\delta \in [0, 1)$. Consider a finite-alphabet channel with $P_{W^n}(y^n | x^n) = P_{Y^n | X^n}(y^n | x^n) = \prod_{i=1}^n P_{Y_i | X_i}(y_i | x_i)$ for all n. For any input X and its corresponding output Y,

$$\underline{I}_{\delta}(\boldsymbol{X};\boldsymbol{Y}) \leq \underline{I}_{\delta}(\bar{\boldsymbol{X}};\bar{\boldsymbol{Y}}) = \underline{I}(\bar{\boldsymbol{X}};\bar{\boldsymbol{Y}}),$$

where \mathbf{Y} is the output due to \mathbf{X} , which is an independent process with the same first order statistics as \mathbf{X} , i.e., $P_{\overline{X}^n}(x^n) = \prod_{i=1}^n P_{X_i}(x_i)$.

System setting:

• Let $\mathcal{X} = \mathcal{Y} = \{0, 1\}$ and

$$Y_i^{(n)} = X_i^{(n)} \oplus Z_i^{(n)}$$

where the arbitrary noise Z is independent of the channel input X.

- Assume that X is a Bernoulli uniform input, i.e., an i.i.d. random process with uniform marginal distribution.
- Then the resultant channel output \boldsymbol{Y} is also Bernoulli uniform no matter what distribution \boldsymbol{Z} has.

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Derivations:

$$\bar{i}(\theta) := \limsup_{n \to \infty} \Pr\left\{\frac{1}{n} \log \frac{P_{Y^n | X^n}(Y^n | X^n)}{P_{Y^n}(Y^n)} \le \theta\right\}$$

$$= \limsup_{n \to \infty} \Pr\left\{\frac{1}{n} \log P_{Z^n}(Z^n) - \frac{1}{n} \log P_{Y^n}(Y^n) \le \theta\right\}$$

$$= \limsup_{n \to \infty} \Pr\left\{\frac{1}{n} \log P_{Z^n}(Z^n) \le \theta - \log(2)\right\}$$

$$= \limsup_{n \to \infty} \Pr\left\{-\frac{1}{n} \log P_{Z^n}(Z^n) \ge \log(2) - \theta\right\}$$

$$= 1 - \liminf_{n \to \infty} \Pr\left\{-\frac{1}{n} \log P_{Z^n}(Z^n) < \log(2) - \theta\right\}.$$

Hence, for $\varepsilon \in (0, 1)$,

$$\begin{split} \underline{I}_{\varepsilon}(\boldsymbol{X};\boldsymbol{Y}) &= \sup\left\{\boldsymbol{\theta}: \overline{i}(\boldsymbol{\theta}) \leq \varepsilon\right\} \\ &= \sup\left\{\boldsymbol{\theta}: 1 - \liminf_{n \to \infty} \Pr\left\{-\frac{1}{n}\log P_{Z^{n}}(Z^{n}) < \log(2) - \boldsymbol{\theta}\right\} \leq \varepsilon\right\} \\ &= \sup\left\{\boldsymbol{\theta}: \liminf_{n \to \infty} \Pr\left\{-\frac{1}{n}\log P_{Z^{n}}(Z^{n}) < \log(2) - \boldsymbol{\theta}\right\} \geq 1 - \varepsilon\right\} \\ &= \sup\left\{(\log(2) - \beta): \liminf_{n \to \infty} \Pr\left\{-\frac{1}{n}\log P_{Z^{n}}(Z^{n}) < \beta\right\} \geq 1 - \varepsilon\right\} \\ &= \log(2) + \sup\left\{-\beta: \liminf_{n \to \infty} \Pr\left\{-\frac{1}{n}\log P_{Z^{n}}(Z^{n}) < \beta\right\} \geq 1 - \varepsilon\right\} \\ &= \log(2) - \inf\left\{\beta: \liminf_{n \to \infty} \Pr\left\{-\frac{1}{n}\log P_{Z^{n}}(Z^{n}) < \beta\right\} \geq 1 - \varepsilon\right\} \\ &= \log(2) - \sup\left\{\beta: \liminf_{n \to \infty} \Pr\left\{-\frac{1}{n}\log P_{Z^{n}}(Z^{n}) < \beta\right\} < 1 - \varepsilon\right\} \\ &= \log(2) - \sup\left\{\beta: \liminf_{n \to \infty} \Pr\left\{-\frac{1}{n}\log P_{Z^{n}}(Z^{n}) < \beta\right\} < 1 - \varepsilon\right\} \\ &= \log(2) - \sup\left\{\beta: \liminf_{n \to \infty} \Pr\left\{-\frac{1}{n}\log P_{Z^{n}}(Z^{n}) < \beta\right\} < 1 - \varepsilon\right\} \\ &= \log(2) - \sup\left\{\beta: \liminf_{n \to \infty} \Pr\left\{-\frac{1}{n}\log P_{Z^{n}}(Z^{n}) < \beta\right\} < 1 - \varepsilon\right\} \\ &= \log(2) - \sup\left\{\beta: \liminf_{n \to \infty} \Pr\left\{-\frac{1}{n}\log P_{Z^{n}}(Z^{n}) \leq \beta\right\} < 1 - \varepsilon\right\} \\ &= \log(2) - \lim_{\delta \uparrow (1 - \varepsilon)} \overline{H}_{\delta}(\boldsymbol{Z}). \end{split}$$

Also, for $\varepsilon \in (0, 1)$,

$$\underline{I}_{\varepsilon}(\boldsymbol{X};\boldsymbol{Y}) \geq \sup\left\{\boldsymbol{\theta}: \limsup_{n \to \infty} \Pr\left[\frac{1}{n}\log\frac{P_{X^{n},Y^{n}}(X^{n},Y^{n})}{P_{X^{n}}(X^{n})P_{Y^{n}}(Y^{n})} < \boldsymbol{\theta}\right] < \varepsilon\right\} \quad (1.5.3)$$

$$= \log(2) - \sup\left\{\boldsymbol{\beta}: \liminf_{n \to \infty} \Pr\left\{-\frac{1}{n}\log P_{Z^{n}}(Z^{n}) \leq \boldsymbol{\beta}\right\} \leq 1 - \varepsilon\right\}$$

$$= \log(2) - \bar{H}_{(1-\varepsilon)}(\boldsymbol{Z}),$$

where (1.5.3) follows from the fact described in Footnote 3. Therefore,

$$\log(2) - \bar{H}_{(1-\varepsilon)}(\boldsymbol{Z}) \leq \underline{I}_{\varepsilon}(\boldsymbol{X};\boldsymbol{Y}) \leq \log(2) - \lim_{\gamma \uparrow (1-\varepsilon)} \bar{H}_{\gamma}(\boldsymbol{Z}) \quad \text{for } \varepsilon \in (0,1).$$

By taking $\varepsilon \downarrow 0$, we obtain

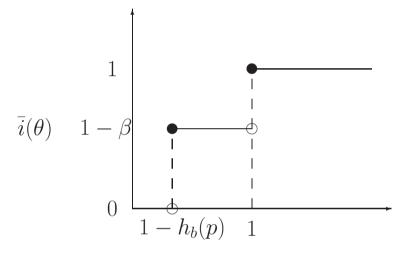
$$\underline{I}(\boldsymbol{X};\boldsymbol{Y}) = \underline{I}_0(\boldsymbol{X};\boldsymbol{Y}) = \log(2) - \bar{H}(\boldsymbol{Z}).$$

Based on this result, we can now compute $I_{\varepsilon}(\mathbf{X}; \mathbf{Y})$ for some specific examples.

Example 1.8

$$\boldsymbol{Z} = \begin{cases} \text{ all-zero sequence with probability } \beta; \\ \text{Bernoulli (with parameter } p) with probability 1 - \beta. \end{cases}$$

Then



Therefore,

$$\underline{I}_{\varepsilon}(\boldsymbol{X};\boldsymbol{Y}) = \begin{cases} 1 - h_b(p), & \text{if } 0 < \varepsilon < 1 - \beta; \\ 1, & \text{if } 1 - \beta \le \varepsilon < 1. \end{cases}$$

Example 1.9 Z =non-stationary binary independent sequence with

$$\Pr\left\{Z_{i}^{(n)}=0\right\}=1-\Pr\left\{Z_{i}^{(n)}=1\right\}=p_{i},$$

then by the fact that

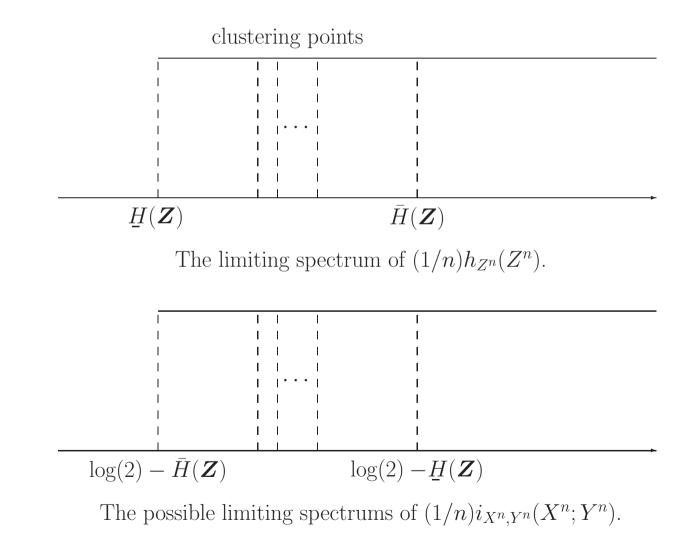
$$\operatorname{Var}\left[-\log P_{Z_{i}^{(n)}}(Z_{i}^{(n)})\right] \leq E\left[\left(\log P_{Z_{i}^{(n)}}(Z_{i}^{(n)})\right)^{2}\right]$$
$$\leq \sup_{0 < p_{i} < 1}\left[p_{i}(\log p_{i})^{2} + (1 - p_{i})(\log(1 - p_{i}))^{2}\right]$$
$$< \log(2),$$

we have (by Chebyshev's inequality) that as $n \to \infty$,

$$\Pr\left\{\left|-\frac{1}{n}\log P_{Z^n}(Z^n) - \frac{1}{n}\sum_{i=1}^n H\left(Z_i^{(n)}\right)\right| > \gamma\right\} \to 0,$$

for any $\gamma > 0$.

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Therefore, $\bar{H}_{(1-\varepsilon)}(\boldsymbol{Z})$ is equal to

$$\bar{H}_{(1-\varepsilon)}(\boldsymbol{Z}) = \begin{cases} \bar{H}(\boldsymbol{Z}) = \limsup_{n \to \infty} \frac{1}{n} \sum_{i=1}^{n} H\left(Z_{i}^{(n)}\right) \\ = \limsup_{n \to \infty} \frac{1}{n} \sum_{i=1}^{n} h_{b}(p_{i}) \\ +\infty, \qquad \text{for } \varepsilon = 0. \end{cases}$$

Consequently,

$$\underline{I}_{\varepsilon}(\boldsymbol{X};\boldsymbol{Y}) = \begin{cases} 1 - \bar{H}(\boldsymbol{Z}) = 1 - \limsup_{n \to \infty} \frac{1}{n} \sum_{i=1}^{n} h_b(p_i), & \text{for } \varepsilon \in [0,1), \\ \infty, & \text{for } \varepsilon = 1. \end{cases}$$