# **Chapter 1**

# **Generalized Information Measures for Arbitrary Systems with Memory**

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### Shannon's entropy II: 1-1

• Entropy of a discrete random variable  $X$ :

$$
H(X) := -\sum_{x \in \mathcal{X}} P_X(x) \log P_X(x) = E_X \left[ -\log P_X(X) \right] \text{ nats}
$$

is a measure of the average amount of uncertainty in  $X$ .

• Entropy rate for a sequence of random variables  $X_1, X_2, \ldots, X_n, \ldots$  is

$$
\lim_{n \to \infty} \frac{1}{n} H(X^n) = \lim_{n \to \infty} \frac{1}{n} E \left[ -\log P_{X^n}(X^n) \right],
$$

assuming the limit exists.

- Operation meaning: Shannon's source coding theorems for stationary and ergodic systems.
- Question: Does these measures have the same operational significance for systems with time-varying and non-stationary statistics. Answer: No.
- Solution: Require new entropy measure which can appropriately characterize the operational limits of *arbitrary* stochastic systems.

### Arbitrary systems with memory II: 1-2

- In general, there are two indices for random variables or observations: <sup>a</sup> time index and <sup>a</sup> space index.
- When a sequence of random variables is denoted by

$$
X_1, X_2, \ldots, X_n, \ldots,
$$

the subscript i of  $X_i$  can be treated as either a time index or a space index, but not both.

• Hence, when <sup>a</sup> sequence of random variables is <sup>a</sup> function of both time and space, the notation of  $X_1, X_2, \ldots, X_n, \ldots$ , is by no means sufficient; and therefore, <sup>a</sup> new model for <sup>a</sup> general time-varying source, such as

$$
X_1^{(n)}, X_2^{(n)}, \ldots, X_t^{(n)}, \ldots,
$$

where  $t$  is the time index and  $n$  is the space or position index (or vice versa), becomes significant.

### Arbitrary systems with memory II: 1-3

• When block-wise (fixed-length) compression of such source (with blocklength <sup>n</sup>) is considered, the same question as to the compression of i.i.d. source arises:

> *what is the minimum compression rate (say in bits per source sample) for which the probability of error probability can be made arbitrarily small as the blocklength goes to infinity?*

• To answer the question, information theorists have to find <sup>a</sup> sequence of data compression codes for each blocklength  $n$  and investigate if the decompression error goes to zero as  $n$  approaches infinity.

### Arbitrary systems with memory II: 1-4

• However, unlike those simple source models such as discrete memorylessness, the source being arbitrary may exhibit distinct statistics for each blocklength  $n; e.g., for$ 

$$
n = 1 : X_1^{(1)}
$$
  
\n
$$
n = 2 : X_1^{(2)}, X_2^{(2)}
$$
  
\n
$$
n = 3 : X_1^{(3)}, X_2^{(3)}, X_3^{(3)}
$$
  
\n
$$
n = 4 : X_1^{(4)}, X_2^{(4)}, X_3^{(4)}, X_4^{(4)}
$$
\n(1.0.2)

the statistics of  $X_1^{(4)}$  could be different from  $X_1^{(1)}$ ,  $X_1^{(2)}$  and  $X_1^{(3)}$  (i.e., the source statistics are not necessarily consistent).

• Since the model in question  $(1.0.2)$  is general, and the system statistics can be *arbitrarily* defined, it is therefore named an *arbitrary system with memory*.

# Arbitrary systems with memory  $\text{II: 1-5}$

• The triangular array of random variables in (1.0.2) is denoted by <sup>a</sup> boldface letter as

$$
\mathbf{X}:=\left\{X^n\right\}_{n=1}^\infty,
$$

where

$$
X^{n}:=\left(X_{1}^{(n)}, X_{2}^{(n)}, \ldots, X_{n}^{(n)}\right);
$$

for convenience, we also write

$$
\mathbf{X} = \left\{ X^n = \left( X_1^{(n)}, X_2^{(n)}, \dots, X_n^{(n)} \right) \right\}_{n=1}^{\infty}.
$$

### $Spectrum and Quantile$  II: 1-6

**Definition 1.1 (Inf/sup-spectrum)** If  $\{A_n\}_{n=1}^{\infty}$  is a sequence of random variables, then its  $\inf\text{-}spectrum\underline{u}(\cdot)$  and its  $sup\text{-}spectrum\bar{u}(\cdot)$  are defined by

$$
\underline{u}(\theta) := \liminf_{n \to \infty} \Pr\{A_n \le \theta\} \quad \text{and} \quad \overline{u}(\theta) := \limsup_{n \to \infty} \Pr\{A_n \le \theta\},\
$$

respectively, where  $\theta \in \mathbb{R}$ .

•  $\underline{u}(\cdot)$  and  $\bar{u}(\cdot)$  are respectively the liminf and the limsup of the cumulative distribution function (CDF) of  $A_n$ .

**Definition 1.2 (Quantile of**  $\inf/\sup$ **-spectrum)** For any  $0 \le \delta \le 1$ , For any  $0 \leq \delta \leq 1$ , the *quantile*  $U_{\delta}$  of the sup-spectrum  $\bar{u}(\cdot)$  and the *quantile*  $\bar{U}_{\delta}$  of the inf-spectrum  $\underline{u}(\cdot)$  are defined by

$$
\underline{U}_{\delta} := \sup \{ \theta : \bar{u}(\theta) \le \delta \} \quad \text{and} \quad \bar{U}_{\delta} := \sup \{ \theta : \underline{u}(\theta) \le \delta \},
$$

respectively. It follows from the above definitions that  $U_{\delta}$  and  $\bar{U}_{\delta}$  are right-continuous and non-decreasing in  $\delta$ . Note that the supremum of an empty set is defined to be −∞.

• If  $\bar{u}(\cdot)$  is strictly increasing, then the quantile is exactly its inverse:  $U_{\delta} = \bar{u}^{-1}(\delta)$ .

Liminf in probability and limsup in probability  $II: 1-7$ 

• *liminf* in probability U of  $\{A_n\}_{n=1}^{\infty}$  is the largest extended real number such that for all  $\xi > 0$ ,

$$
\lim_{n \to \infty} \Pr[A_n \le U - \xi] = 0.
$$

• *limsup* in probability  $\bar{U}$  of  $\{A_n\}_{n=1}^{\infty}$  is the smallest extended real number such that for all  $\xi > 0$ ,

 $\lim_{n\to\infty} \Pr[A_n \ge \bar{U} + \xi] = 0.$ 

$$
\underline{U} = \lim_{\delta \downarrow 0} \underline{U}_{\delta} = \underline{U}_0
$$

and

•

$$
\bar{U} = \lim_{\delta \uparrow 1} \bar{U}_{\delta} = \sup \{ \theta : \underline{u}(\theta) < 1 \}.
$$

• It readily follows from the above definitions that

$$
\underline{U} \le \underline{U}_{\delta} \le \overline{U}_{\delta} \le \overline{U} \quad \text{for } \delta \in [0, 1).
$$

 $\overline{U}_1 = \underline{U}_1 = \infty.$ 



The asymptotic CDFs (spectrums) of  $\{A_n\}_{n=1}^{\infty}$  and their quantiles:  $\bar{u}(\cdot) = \text{sup-spectrum of }\{A_n\}, \quad \underline{u}(\cdot) = \text{inf-spectrum of }\{A_n\},$  $U_{\delta} =$  quantile of  $\bar{u}(\cdot),$   $\qquad \qquad \bar{U}_{\delta} =$  quantile of  $\underline{u}(\cdot),$  $\underline{U}=\lim_{\delta\downarrow 0}\underline{U}_\delta=\underline{U}_0,\qquad \underline{U}_{1^-}=\lim_{\xi\uparrow 1}\underline{U}_\xi,\qquad \bar{U}=\lim_{\delta\uparrow 1}\bar{U}_\delta.$ 

### Properties of quantile II: 1-9

#### **Lemma 1.4** Assume:

- Two random sequences:  $\{A_n\}_{n=1}^{\infty}$  and  $\{B_n\}_{n=1}^{\infty}$ ;
- $\bar{u}(\cdot)$  = sup-spectrum of  $\{A_n\}_{n=1}^{\infty}$ ;  $U_{\delta}$  = quantile of  $\bar{u}(\cdot)$ ;
- $\underline{u}(\cdot) = \inf\text{-spectrum of }\{A_n\}_{n=1}^{\infty}$ ;  $\overline{U}_d$  $\zeta_{\delta} =$  quantile of  $\underline{u}(\cdot);$
- $\bar{v}(\cdot)$  = sup-spectrum of  $\{B_n\}_{n=1}^{\infty}$ ;  $V_\delta$  = quantile of  $\bar{u}(\cdot)$ ;
- $\underline{v}(\cdot) = \inf\{\operatorname{spectrum of}\{B_n\}_{n=1}^{\infty}; \overline{V}_{\delta}$  $V_{\delta} =$  quantile of  $\underline{u}(\cdot);$

• 
$$
(\overline{u+v})(\cdot)
$$
 = sup-spectrum of  $\{A_n+B_n\}_{n=1}^{\infty}$ , i.e.,

$$
(\overline{u+v})(\theta) := \limsup_{n \to \infty} \Pr\{A_n + B_n \le \theta\};
$$

 $(\underline{U} + V)_{\delta} =$  quantile of  $(\overline{u+v})(\cdot);$ 

\n- \n
$$
\begin{aligned}\n \Phi\left(\underline{u} + \underline{v}\right)(\cdot) &= \inf\text{-spectrum of }\{A_n + B_n\}_{n=1}^{\infty}, \text{ i.e.,} \\
 &\quad (\underline{u} + \underline{v})(\theta) := \liminf_{n \to \infty} \Pr\{A_n + B_n \leq \theta\}; \\
 &\quad (\overline{U} + \overline{V})_{\delta} = \text{quantile of } (\underline{u} + \underline{v})(\cdot).\n \end{aligned}
$$
\n
\n

## Properties of quantile II: 1-10

Then the following statements hold.

- 1.  $U_{\delta}$  and  $\bar{U}_{\delta}$  $U_{\delta}$  are both non-decreasing and right-continuous functions of  $\delta$  for  $\delta \in [0, 1].$
- 2.  $\lim_{\delta \downarrow 0} U_{\delta} = U_0$  and  $\lim_{\delta \downarrow 0} \bar{U}_{\delta}$  $\bar U_\delta = \bar U_0$  $\cup$   $0$  .
- 3. For  $\delta \geq 0, \gamma \geq 0$ , and  $\delta + \gamma \leq 1$ ,

$$
(\underline{U+V})_{\delta+\gamma} \ge \underline{U}_{\delta} + \underline{V}_{\gamma},\tag{1.2.1}
$$

and

$$
(\overline{U+V})_{\delta+\gamma} \ge \underline{U}_{\delta} + \bar{V}_{\gamma}.
$$
\n(1.2.2)

4. For  $\delta \geq 0, \gamma \geq 0$ , and  $\delta + \gamma \leq 1$ ,

$$
(\underline{U+V})_{\delta} \le \underline{U}_{\delta+\gamma} + \bar{V}_{(1-\gamma)},\tag{1.2.3}
$$

and

$$
(\overline{U+V})_{\delta} \le \overline{U}_{\delta+\gamma} + \overline{V}_{(1-\gamma)}.
$$
\n(1.2.4)

### Generalized information measures II: 1-11

In Definitions 1.1 and 1.2,

 $\bullet$   $A_n$  = *normalized entropy density*, i.e.,

$$
\frac{1}{n}h_{X^n}(X^n) := -\frac{1}{n}\log P_{X^n}(X^n),
$$

 $\delta$ *-inf-entropy*  $\textit{rate}_{\mathcal{H}}(X) =$  quantile of the sup-spectrum of 1  $\, n$  $h_{X^n}(X^n)$ 

 $\delta$ -*sup-entropy* rate  $\bar{H}_{\delta}(\boldsymbol{X}) =$  quantile of the inf-spectrum of 1  $\, n$  $h_{X^n}(X^n).$ 

• 
$$
A_n
$$
 = normalized information density, i.e.,

$$
\frac{1}{n}i_{X^nW^n}(X^n;Y^n) = \frac{1}{n}i_{X^n,Y^n}(X^n;Y^n) := \frac{1}{n}\log\frac{P_{X^n,Y^n}(X^n,Y^n)}{P_{X^n}(X^n)P_{Y^n}(Y^n)},
$$

 $\delta$ -*inf-information-rate*  $I_{\delta}(\boldsymbol{X}; \boldsymbol{Y}) =$  quantile of the sup-spectrum of 1  $\, n$  $i_{X^nW^n}(X^n;Y^n)$ 

 $\delta$ -*sup-information-rate*  $\bar{I}_{\delta}(\boldsymbol{X}; \boldsymbol{Y}) =$  quantile of the inf-spectrum of 1  $\, n$  $i_{X^nW^n}(X^n;Y^n).$ 

### Generalized information measures II: 1-12

 $\bullet$   $A_n$  = *normalized log-likelihood ratio*, i.e.,

$$
\frac{1}{n}d_{X^n}(X^n||\hat{X}^n) := \frac{1}{n}\log\frac{P_{X^n}(X^n)}{P_{\hat{X}^n}(X^n)}
$$

δ-*inf-divergence rate*  $D_{\delta}(\boldsymbol{X}||\hat{\boldsymbol{X}})$  = quantile of the sup-spectrum of  $\frac{1}{\epsilon}$  $\, n$  $d_{X^n}(X^n\|\hat{X}^n)$ 

δ-*sup-divergence rate*  $\bar{D}_δ(\boldsymbol{X}||\hat{\boldsymbol{X}})$  = quantile of the inf-spectrum of  $\frac{1}{\epsilon}$  $\, n$  $d_{X^n}(X^n\|\hat{X}^n).$ 

- Notes:
	- $-$  The *inf-entropy-rate*  $\underline{H}(\boldsymbol{X})$  and the *sup-entropy-rate*  $\overline{H}(\boldsymbol{X})$  are special cases of the  $\delta$ -inf/sup-entropy rate measures:

$$
\underline{H}(\mathbf{X}) = \underline{H}_0(\mathbf{X}), \text{ and } \overline{H}(\mathbf{X}) = \lim_{\delta \uparrow 1} \overline{H}_\delta(\mathbf{X}).
$$

 $\sim$  Concept: If the random variable  $(1/n)h(X^n)$  exhibits a limiting distribution, and suppose the limiting distribution of  $(1/n)h_{X^n}(X^n)$  is positive over  $(-2, 2)$ ; and zero, otherwise. Then  $\bar{H}(\boldsymbol{X}) = 2$  and  $H(\boldsymbol{X}) = -2$ .

# Generalized information measures II: 1-13



Generalized entropy measures where  $\delta \in [0,1]$ .

# Generalized information measures





Generalized mutual information measures where  $\delta \in [0,1]$ .

# Generalized information measures





Generalized divergence measures where  $\delta \in [0,1]$ .

• An example of basic properties for Shannon's entropy:  $I(X;Y) = H(Y) - H(Y|X).$ 

 $-$  By taking  $\delta = 0$  and letting  $\gamma \downarrow 0$  in

 $(\underline{U} + V)_{\delta + \gamma} \ge \underline{U}_{\delta} + \underline{V}_{\gamma}$  for  $\delta \ge 0, \gamma \ge 0$ , and  $\delta + \gamma \le 1$ 

and

$$
(\underline{U+V})_{\delta} \le \underline{U}_{\delta+\gamma} + \bar{V}_{(1-\gamma)} \text{ for } \delta \ge 0, \gamma \ge 0, \text{ and } \delta+\gamma \le 1,
$$

we obtain

$$
(\underline{U+V}) \ge \underline{U}_0 + \lim_{\gamma \downarrow 0} \underline{V}_\gamma \ge \underline{U} + \underline{V}
$$

and

$$
(\underline{U+V}) \le \lim_{\gamma \downarrow 0} \underline{U}_{\gamma} + \lim_{\gamma \downarrow 0} \overline{V}_{(1-\gamma)} = \underline{U} + \overline{V}.
$$

 $-$  Meaning: The liminf in probability of a sequence of random variables  $A_n +$  $B_n$  is upper bounded by the liminf in probability of  $A_n$  plus the limsup in probability of  $B_n$ ; and is lower bounded by the sum of the liminfs in probability of  $A_n$  and  $B_n$ .

**–** This fact can be used to show that

$$
I(\boldsymbol{X};\boldsymbol{Y}) + \underline{H}(\boldsymbol{Y}|\boldsymbol{X}) \le \underline{H}(\boldsymbol{Y}) \le \underline{I}(\boldsymbol{X};\boldsymbol{Y}) + \bar{H}(\boldsymbol{Y}|\boldsymbol{X}),
$$

or equivalently,

$$
H(Y) - \bar{H}(Y|X) \leq I(X;Y) \leq H(Y) - H(Y|X).
$$

**Lemma 1.5** For a source  $X$  with finite alphabet  $\mathcal{X}$  and arbitrary sources  $Y$  and *Z*, the following properties hold.

- 1.  $\bar{H}_{\delta}(\boldsymbol{X}) \geq 0$  for  $\delta \in [0,1]$ . (This property also applies to  $H_{\delta}(\boldsymbol{X}), \bar{I}_{\delta}(\boldsymbol{X}; \boldsymbol{Y}),$  $I_{\delta}(\boldsymbol{X};\boldsymbol{Y}),\,\bar{D}_{\delta}(\boldsymbol{X}\|\hat{\boldsymbol{X}}),\,\text{and}\,\underline{D}_{\delta}(\boldsymbol{X}\|\hat{\boldsymbol{X}}).$
- 2.  $I_{\delta}(\boldsymbol{X}; \boldsymbol{Y}) = I_{\delta}(\boldsymbol{Y}; \boldsymbol{X})$  and  $\bar{I}_{\delta}(\boldsymbol{X}; \boldsymbol{Y}) = \bar{I}_{\delta}(\boldsymbol{Y}; \boldsymbol{X})$  for  $\delta \in [0, 1]$ .

3. For  $0 \leq \delta < 1, 0 \leq \gamma < 1$  and  $\delta + \gamma \leq 1$ ,

$$
\underline{I}_{\delta}(\boldsymbol{X};\boldsymbol{Y}) \le \underline{H}_{\delta+\gamma}(\boldsymbol{Y}) - \underline{H}_{\gamma}(\boldsymbol{Y}|\boldsymbol{X}), \tag{1.4.1}
$$

$$
I_{\delta}(\boldsymbol{X}; \boldsymbol{Y}) \leq \bar{H}_{\delta + \gamma}(\boldsymbol{Y}) - \bar{H}_{\gamma}(\boldsymbol{Y}|\boldsymbol{X}), \qquad (1.4.2)
$$

$$
\bar{I}_{\delta}(\boldsymbol{X}; \boldsymbol{Y}) \leq \bar{H}_{\delta + \gamma}(\boldsymbol{Y}) - \underline{H}_{\gamma}(\boldsymbol{Y}|\boldsymbol{X}), \tag{1.4.3}
$$

$$
I_{\delta+\gamma}(\boldsymbol{X};\boldsymbol{Y}) \geq H_{\delta}(\boldsymbol{Y}) - \bar{H}_{(1-\gamma)}(\boldsymbol{Y}|\boldsymbol{X}), \qquad (1.4.4)
$$

and

$$
\bar{I}_{\delta+\gamma}(\boldsymbol{X};\boldsymbol{Y}) \ge \bar{H}_{\delta}(\boldsymbol{Y}) - \bar{H}_{(1-\gamma)}(\boldsymbol{Y}|\boldsymbol{X}). \tag{1.4.5}
$$

(Note that the case of  $(\delta, \gamma) = (1, 0)$  holds for  $(1.4.1)$  and  $(1.4.2)$ , and the case of  $(\delta, \gamma) = (0, 1)$  holds for  $(1.4.3), (1.4.4)$  and  $(1.4.5).$ 

4.  $0 \leq H_\delta(\boldsymbol{X}) \leq \bar{H}_\delta(\boldsymbol{X}) \leq \log |\mathcal{X}|$  for  $\delta \in [0,1)$ , where each  $X_i^{(n)}$  takes values in X for  $i = 1, \ldots, n$  and  $n = 1, 2, \ldots$ 

5.  $I_{\delta}(\boldsymbol{X}, \boldsymbol{Y}; \boldsymbol{Z}) \geq I_{\delta}(\boldsymbol{X}; \boldsymbol{Z})$  for  $\delta \in [0, 1]$ .

**Property 1:**

$$
\Pr\left\{-\frac{1}{n}\log P_{X^n}(X^n) < 0\right\} = 0,
$$
\n
$$
\Pr\left\{\frac{1}{n}\log \frac{P_{X^n}(X^n)}{P_{\hat{X}^n}(X^n)} < -\nu\right\} = P_{X^n}\left\{x^n \in \mathcal{X}^n : \frac{1}{n}\log \frac{P_{X^n}(x^n)}{P_{\hat{X}^n}(x^n)} < -\nu\right\}
$$
\n
$$
= \sum_{\{x^n \in \mathcal{X}^n : P_{X^n}(x^n) < P_{\hat{X}^n}(x^n)e^{-nv}\}} P_{\hat{X}^n}(x^n)
$$
\n
$$
\leq \sum_{\{x^n \in \mathcal{X}^n : P_{X^n}(x^n) < P_{\hat{X}^n}(x^n)e^{nv}\}} P_{\hat{X}^n}(x^n)e^{-nv}
$$
\n
$$
\leq e^{-nv} \cdot \sum_{\{x^n \in \mathcal{X}^n : P_{X^n}(x^n) < P_{\hat{X}^n}(x^n)e^{nv}\}} P_{\hat{X}^n}(x^n)
$$
\n
$$
\leq e^{-\nu n}, \qquad (1.4.6)
$$

and, by following the same procedure as (1.4.6),

$$
\Pr\left\{\frac{1}{n}\log\frac{P_{X^n,Y^n}(X^n,Y^n)}{P_{X^n}(X^n)P_{Y^n}(Y^n)}<-\nu\right\} \leq e^{-\nu n}.
$$

**Property 2:** An immediate consequence of the definition.

**Property 3:** Follow from the facts that

$$
\frac{1}{n}h_{Y^n}(Y^n) = \frac{1}{n}i_{X^n,Y^n}(X^n;Y^n) + \frac{1}{n}h_{X^n,Y^n}(Y^n|X^n),
$$

where

$$
\frac{1}{n}h_{X^n,Y^n}(Y^n|X^n) := -\frac{1}{n}\log P_{Y^n|X^n}(Y^n|X^n).
$$

**Property** 4:  $\bar{H}_{\delta}(\cdot)$  is non-decreasing in  $\delta$ ,  $\bar{H}_{\delta}(\boldsymbol{X}) \leq \bar{H}(\boldsymbol{X})$ , and  $\bar{H}(\boldsymbol{X}) \leq$  $log |\mathcal{X}|$ . The last inequality can be proved as follows:

$$
\Pr\left\{\frac{1}{n}h_{X^n}(X^n) \le \log |\mathcal{X}| + \nu\right\} |
$$
  
= 1 - P\_{X^n}\left\{x^n \in \mathcal{X}^n : \frac{1}{n}\log \frac{P\_{X^n}(X^n)}{1/|\mathcal{X}|^n} < -\nu\right\}  
\ge 1 - e^{-n\nu},

where the last step can be obtained by using the same procedure as  $(1.4.6)$ . Therefore,  $h(\log |\mathcal{X}| + \nu) = 1$  for any  $\nu > 0$ , which indicates that  $\bar{H}(\mathbf{X}) \leq$  $\log |\mathcal{X}|$ .

**Property 5:**

$$
\frac{1}{n}i_{X^n,Y^n,Z^n}(X^n,Y^n;Z^n) = \frac{1}{n}i_{X^n,Z^n}(X^n;Z^n) + \frac{1}{n}i_{X^n,Y^n,Z^n}(Y^n;Z^n|X^n).
$$

 $\Box$ 

**Lemma 1.6 (Data processing lemma)** Fix  $\delta \in [0,1]$ . Consider a channel with finite input and output alphabets,  $\mathcal X$  and  $\mathcal Y$ , respectively, and with distribution  $P_{W^n}(y^n|x^n) = P_{Y^n|X^n}(y^n|x^n) = \prod_{i=1}^n Y_i$  $\sum_{i=1}^n P_{Y_i|X_i}(y_i|x_i)$  for all  $n, x^n \in \mathcal{X}^n$  and  $y^n \in \mathcal{Y}^n$ . Then

$$
I_{\delta}(\boldsymbol{X}_1; \boldsymbol{X}_3) \leq I_{\delta}(\boldsymbol{X}_1; \boldsymbol{X}_2).
$$

**Proof:** By Property 5 of Lemma 1.5, we get

$$
I_{\delta}(\boldsymbol{X}_1;\boldsymbol{X}_3)\leq I_{\delta}(\boldsymbol{X}_1;\boldsymbol{X}_2,\boldsymbol{X}_3)=I_{\delta}(\boldsymbol{X}_1;\boldsymbol{X}_2),
$$

where the equality holds because

$$
\frac{1}{n}\log\frac{P_{X_1^n,X_2^n,X_3^n}(x_1^n,x_2^n,x_3^n)}{P_{X_1^n}(x_1^n)P_{X_2^n,X_3^n}(x_2^n,x_3^n)}=\frac{1}{n}\log\frac{P_{X_1^n,X_2^n}(x_1^n,x_2^n)}{P_{X_1^n}(x_1^n)P_{X_2^n}(x_2^n)}.
$$

 $\Box$ 

**Lemma 1.7** (Optimality of independent inputs) Fix  $\delta \in [0,1)$ . Consider a finite-alphabet channel with  $P_{W^n}(y^n|x^n) = P_{Y^n|X^n}(y^n|x^n) = \prod_{i=1}^n P_{X^n}(y^n|x^n)$  $\sum\limits_{i=1}^n P_{Y_i|X_i}(y_i|x_i)$ for all <sup>n</sup>. For any input *X* and its corresponding output *Y* ,

$$
I_{\delta}(\boldsymbol{X};\boldsymbol{Y}) \leq I_{\delta}(\bar{\boldsymbol{X}};\bar{\boldsymbol{Y}}) = I(\bar{\boldsymbol{X}};\bar{\boldsymbol{Y}}),
$$

where  $\bar{Y}$  is the output due to  $\bar{X}$ , which is an independent process with the same first order statistics as  $\boldsymbol{X}$ , i.e.,  $P_{\overline{X}^n}(x^n) = \prod_{i=1}^n$  $\sum_{i=1}^{n} P_{X_i}(x_i)$ .

### Computation of general information measures  $\frac{II!}{II! \cdot 1-23}$

**System setting**:

• Let  $\mathcal{X} = \mathcal{Y} = \{0, 1\}$  and

$$
Y_i^{(n)} = X_i^{(n)} \oplus Z_i^{(n)}
$$

where the arbitrary noise *Z* is independent of the channel input *X*.

- Assume that *X* is <sup>a</sup> Bernoulli uniform input, i.e., an i.i.d. random process with uniform marginal distribution.
- Then the resultant channel output *Y* is also Bernoulli uniform no matter what distribution *Z* has.

# Computation of general information measures  $\qquad$  II: 1-24

Derivations:

$$
\overline{i}(\theta) := \limsup_{n \to \infty} \Pr\left\{ \frac{1}{n} \log \frac{P_{Y^n|X^n}(Y^n|X^n)}{P_{Y^n}(Y^n)} \le \theta \right\}
$$
  
\n
$$
= \limsup_{n \to \infty} \Pr\left\{ \frac{1}{n} \log P_{Z^n}(Z^n) - \frac{1}{n} \log P_{Y^n}(Y^n) \le \theta \right\}
$$
  
\n
$$
= \limsup_{n \to \infty} \Pr\left\{ \frac{1}{n} \log P_{Z^n}(Z^n) \le \theta - \log(2) \right\}
$$
  
\n
$$
= \limsup_{n \to \infty} \Pr\left\{ -\frac{1}{n} \log P_{Z^n}(Z^n) \ge \log(2) - \theta \right\}
$$
  
\n
$$
= 1 - \liminf_{n \to \infty} \Pr\left\{ -\frac{1}{n} \log P_{Z^n}(Z^n) < \log(2) - \theta \right\}.
$$

# Computation of general information measures  $\frac{II!}{II! \cdot 1 \cdot 25}$

Hence, for  $\varepsilon \in (0,1)$ ,

$$
J_{\varepsilon}(\mathbf{X}; \mathbf{Y}) = \sup \{ \theta : \bar{i}(\theta) \le \varepsilon \}
$$
  
\n
$$
= \sup \left\{ \theta : 1 - \liminf_{n \to \infty} \Pr \left\{ -\frac{1}{n} \log P_{Z^n}(Z^n) < \log(2) - \theta \right\} \le \varepsilon \right\}
$$
  
\n
$$
= \sup \left\{ \theta : \liminf_{n \to \infty} \Pr \left\{ -\frac{1}{n} \log P_{Z^n}(Z^n) < \log(2) - \theta \right\} \ge 1 - \varepsilon \right\}
$$
  
\n
$$
= \sup \left\{ (\log(2) - \beta) : \liminf_{n \to \infty} \Pr \left\{ -\frac{1}{n} \log P_{Z^n}(Z^n) < \beta \right\} \ge 1 - \varepsilon \right\}
$$
  
\n
$$
= \log(2) + \sup \left\{ -\beta : \liminf_{n \to \infty} \Pr \left\{ -\frac{1}{n} \log P_{Z^n}(Z^n) < \beta \right\} \ge 1 - \varepsilon \right\}
$$
  
\n
$$
= \log(2) - \inf \left\{ \beta : \liminf_{n \to \infty} \Pr \left\{ -\frac{1}{n} \log P_{Z^n}(Z^n) < \beta \right\} \ge 1 - \varepsilon \right\}
$$
  
\n
$$
= \log(2) - \sup \left\{ \beta : \liminf_{n \to \infty} \Pr \left\{ -\frac{1}{n} \log P_{Z^n}(Z^n) < \beta \right\} < 1 - \varepsilon \right\}
$$
  
\n
$$
\le \log(2) - \sup \left\{ \beta : \liminf_{n \to \infty} \Pr \left\{ -\frac{1}{n} \log P_{Z^n}(Z^n) \le \beta \right\} < 1 - \varepsilon \right\}
$$
  
\n
$$
= \log(2) - \lim_{\delta \uparrow (1-\varepsilon)} \bar{H}_{\delta}(Z).
$$

## Computation of general information measures  $\frac{II!}{II! \cdot 1-26}$

Also, for  $\varepsilon \in (0,1)$ ,

$$
\underline{I}_{\varepsilon}(\boldsymbol{X}; \boldsymbol{Y}) \geq \sup \left\{ \theta : \limsup_{n \to \infty} \Pr \left[ \frac{1}{n} \log \frac{P_{X^n, Y^n}(X^n, Y^n)}{P_{X^n}(X^n) P_{Y^n}(Y^n)} < \theta \right] < \varepsilon \right\} \quad (1.5.3)
$$
  
=  $\log(2) - \sup \left\{ \beta : \liminf_{n \to \infty} \Pr \left\{ -\frac{1}{n} \log P_{Z^n}(Z^n) \leq \beta \right\} \leq 1 - \varepsilon \right\}$   
=  $\log(2) - \bar{H}_{(1-\varepsilon)}(\boldsymbol{Z}),$  (1.5.3)

where (1.5.3) follows from the fact described in Footnote 3. Therefore,

$$
\log(2) - \bar{H}_{(1-\varepsilon)}(\mathbf{Z}) \le \underline{I}_{\varepsilon}(\mathbf{X}; \mathbf{Y}) \le \log(2) - \lim_{\gamma \uparrow (1-\varepsilon)} \bar{H}_{\gamma}(\mathbf{Z}) \quad \text{for } \varepsilon \in (0,1).
$$

By taking  $\varepsilon \downarrow 0$ , we obtain

$$
I(\boldsymbol{X};\boldsymbol{Y})=I_0(\boldsymbol{X};\boldsymbol{Y})=\log(2)-\bar{H}(\boldsymbol{Z}).
$$

Based on this result, we can now compute  $I_{\varepsilon}(\bm{X};\bm{Y})$  for some specific examples.

# Computation of general information measures  $\frac{II!}{II! \cdot 1-27}$

**Example 1.8**

$$
\mathbf{Z} = \begin{cases} \text{ all-zero sequence with probability } \beta; \\ \text{Bernoulli (with parameter } p) \text{ with probability } 1 - \beta. \end{cases}
$$

Then

$$
\frac{1}{n}h_{Z^n}(Z^n) \to \begin{cases} 0, & \text{with probability } \beta; \\ h_b(p), & \text{with probability } 1 - \beta, \end{cases}
$$
  
where  $h_b(p) := -p \log p - (1 - p) \log(1 - p)$ .  
  

$$
\frac{1}{\log p}
$$

### Computation of general information measures  $\frac{II!}{II! \cdot 1-28}$



Therefore,

$$
I_{\varepsilon}(\mathbf{X}; \mathbf{Y}) = \begin{cases} 1 - h_b(p), & \text{if } 0 < \varepsilon < 1 - \beta; \\ 1, & \text{if } 1 - \beta \le \varepsilon < 1. \end{cases}
$$

**Example 1.9** *Z* =non-stationary binary independent sequence with

$$
\Pr\left\{Z_i^{(n)} = 0\right\} = 1 - \Pr\left\{Z_i^{(n)} = 1\right\} = p_i,
$$

then by the fact that

$$
\operatorname{Var}\left[-\log P_{Z_i^{(n)}}(Z_i^{(n)})\right] \le E\left[\left(\log P_{Z_i^{(n)}}(Z_i^{(n)})\right)^2\right] \\
\le \sup_{0 < p_i < 1} \left[p_i(\log p_i)^2 + (1 - p_i)(\log(1 - p_i))^2\right] \\
&< \log(2),
$$

we have (by Chebyshev's inequality) that as  $n \to \infty$ ,

$$
\Pr\left\{ \left| -\frac{1}{n} \log P_{Z^n}(Z^n) - \frac{1}{n} \sum_{i=1}^n H\left(Z_i^{(n)}\right) \right| > \gamma \right\} \to 0,
$$

for any  $\gamma > 0$ .

## Computation of general information measures  $\frac{II!}{II! \cdot 1-30}$



# Computation of general information measures  $\mathbb{I}$ : 1-31

Therefore,  $\bar{H}_{(1-\varepsilon)}(\boldsymbol{Z})$  is equal to

$$
\bar{H}_{(1-\varepsilon)}(\mathbf{Z}) = \begin{cases}\n\bar{H}(\mathbf{Z}) = \limsup_{n \to \infty} \frac{1}{n} \sum_{i=1}^{n} H\left(Z_i^{(n)}\right) \\
= \limsup_{n \to \infty} \frac{1}{n} \sum_{i=1}^{n} h_b(p_i) \\
+\infty, \quad \text{for } \varepsilon = 0.\n\end{cases}
$$

Consequently,

$$
\underline{I}_{\varepsilon}(\mathbf{X}; \mathbf{Y}) = \begin{cases} 1 - \bar{H}(\mathbf{Z}) = 1 - \limsup_{n \to \infty} \frac{1}{n} \sum_{i=1}^{n} h_b(p_i), & \text{for } \varepsilon \in [0, 1), \\ \infty, & \text{for } \varepsilon = 1. \end{cases}
$$