1. (50%) Prove that among all continuous random variables of mean  $\mu$  and variance  $\sigma^2$ , Gaussian random variable has the largest differential entropy.

Hint: Let X be the Gaussian random variable of mean  $\mu$  and variance  $\sigma^2$ . Let Y be an arbitrary random variable of the same mean and variance as X. Then,

$$h(X) = \int_{\Re} f_X(x) \log_2 \frac{1}{f_X(x)} dx$$
  

$$= \int_{\Re} f_X(x) \left[ \frac{1}{2} \log_2(2\pi\sigma^2) + \log_2(e) \cdot \frac{(x-\mu)^2}{2\sigma^2} \right] dx$$
  

$$= \int_{\Re} f_Y(x) \left[ \frac{1}{2} \log_2(2\pi\sigma^2) + \log_2(e) \cdot \frac{(x-\mu)^2}{2\sigma^2} \right] dx$$
  

$$= \int_{\Re} f_Y(x) \log_2 \frac{1}{f_X(x)} dx$$
(1)

Complete the proof based on (1).

**Solution.** See Slide IDC 6-54. Other than the fundamental inequality, you can also use Jensen's inequality, or log-sum inequality (or the non-negativity of divergence) to prove the result.

2. (50%) Let the generator polynomial of a polynomial code of length n = 5 be  $g(X) = X^3 + X + 1$ . List all the code polynomials of this polynomial code via

$$c(X) = a(X)g(X).$$

What is the minimum pairwise Hamming distance of this code? Hint:

$$d_{\min} = \min_{\boldsymbol{c}_i, \boldsymbol{c}_j \in \mathcal{C}, \boldsymbol{c}_i \neq \boldsymbol{c}_j} d_{\mathrm{H}}(\boldsymbol{c}_i, \boldsymbol{c}_j) = \min_{\boldsymbol{c} \in \mathcal{C}, \boldsymbol{c} \neq \boldsymbol{0}} w_{\mathrm{H}}(\boldsymbol{c})$$

**Solution.**  $c(X) = (a_0 + a_1 X)(1 + X + X^3)$  implies

c(X)	$(a_0 + a_1 X)(1 + X + X^3)$
$0 + 0 \cdot X + 0 \cdot X^2 + 0 \cdot X^3 + 0 \cdot X^4$	$(0+0\cdot X)(1+X+X^3)$
$0 + 1 \cdot X + 1 \cdot X^2 + 0 \cdot X^3 + 1 \cdot X^4$	$(0+1\cdot X)(1+X+X^3)$
$1 + 1 \cdot X + 0 \cdot X^2 + 1 \cdot X^3 + 0 \cdot X^4$	$(1+0\cdot X)(1+X+X^3)$
$1+0\cdot X+1\cdot X^2+1\cdot X^3+1\cdot X^4$	$(1+1\cdot X)(1+X+X^3)$

The minimum pairwise Hamming distance is the smallest Hamming weight of non-zero codewords, which is 3.