## Digital Communications Chapter 9 Digital Communications Through Band-Limited Channels

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# 9.2 Signal design for band-limited channels

For the baseband waveform

$$s_{\ell}(t) = \sum_{n=-\infty}^{\infty} I_n g(t - nT)$$

- Channel is band-limited to [-W, W].
- Transmitted signals outside [-W, W] will be truncated.

How to design g(t) to yield optimal performance?

<sup>1</sup> Here we use  $s_{\ell}(t)$  instead of v(t) as being used in text to clearly indicate that  $s_{\ell}(t)$  is the lowpass equivalent signal of s(t).

Recall that

$$s_{\ell}(t) = \sum_{n=-\infty}^{\infty} I_n g(t-nT)$$

is a cyclostationary random process with i.i.d. discrete random process  $\{I_n\}$ .

• Autocorrelation function of  $s_{\ell}(t)$  is

$$R_{s_{\ell}}(t+\tau,t) = \mathbb{E}[s_{\ell}(t+\tau)s_{\ell}^{*}(t)]$$

is periodic with period T.

• The time-average autocorrelation function

$$\bar{R}_{s_{\ell}}(\tau) = \frac{1}{T} \int_{-\frac{T}{2}}^{\frac{T}{2}} R_{s_{\ell}}(t+\tau,t) dt$$
$$= \frac{1}{T} \sum_{n=-\infty}^{\infty} R_{I}(n) \int_{-\infty}^{\infty} g(t+\tau-nT)g^{*}(t) dt$$

$$\bar{R}_{s_{\ell}}(\tau) = \frac{1}{T} \sum_{n} R_{I}(n) \int_{-\infty}^{\infty} g(t + \tau - nT) g^{*}(t) dt$$

• The time-average power spectral density of  $s_{\ell}(t)$  is

$$\bar{S}_{s_{\ell}}(f) = \int_{-\infty}^{\infty} \bar{R}_{s_{\ell}}(\tau) e^{-i2\pi f\tau} d\tau$$
$$= \frac{1}{T} \left[ \sum_{n=-\infty}^{\infty} R_{I}(n) e^{-i2\pi fnT} \right] |G(f)|^{2}$$

• Assuming  $I_n$  is zero mean and i.i.d.,  $R_I(n) = \sigma^2 \delta_n$ , hence

$$\bar{S}_{s_{\ell}}(f) = \frac{\sigma^2}{T} |G(f)|^2$$

For static channel with impulse response  $c_{\ell}(t)$ , band-limited to [-W, W], i.e.

$$C_\ell(f) = 0$$
 for  $|f| > W$ 

The received signal is

$$r_{\ell}(t) = c_{\ell}(t) \star s_{\ell}(t) + n_{\ell}(t)$$
  
=  $\int_{-\infty}^{\infty} c_{\ell}(\tau) \left( \sum_{n=-\infty}^{\infty} I_n g(t - \tau - nT) \right) d\tau + n_{\ell}(t)$   
=  $\sum_{n=-\infty}^{\infty} I_n \int_{-\infty}^{\infty} g(t - nT - \tau) c_{\ell}(\tau) d\tau + n_{\ell}(t)$ 

whose time-average power spectral density is

$$\bar{S}_{r_{\ell}}(f) = \frac{\sigma^2}{T} |G(f)|^2 |C_{\ell}(f)|^2 + 2N_0 \operatorname{rect}\left(\frac{f}{2W}\right)$$

Note that  $n_{\ell}(t)$  is a band-limited white noise process.

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$$r_{\ell}(t) = \sum_{n=-\infty}^{\infty} I_n h_{\ell}(t-nT) + n_{\ell}(t) \text{ with } h_{\ell}(t) = g(t) \star c_{\ell}(t)$$

- Assume channel Impulse response  $c_{\ell}(t)$  known to Rx.
- The match-filtered received signal is

$$y_{\ell}(t) = r_{\ell}(t) \star h_{\ell}^{*}(T-t) = \sum_{n} I_{n} x_{\ell}(t-nT) + z_{\ell}(t)$$
  
where  $x_{\ell}(t) = h_{\ell}(t) \star h_{\ell}^{*}(T-t)$  and  
 $z_{\ell}(t) = n_{\ell}(t) \star h_{\ell}^{*}(T-t).$   
For simplicity, one may use  $h_{\ell}^{*}(-t)$  instead of  $h_{\ell}^{*}(T-t)$ 

• Sampling at t = kT,  $k \in \mathbb{Z}$ , we get

$$y_{k} = y_{\ell}(kT) = \sum_{n=-\infty}^{\infty} I_{n} x_{\ell}(kT - nT) + z_{\ell}(kT) = \sum_{n} I_{n} x_{k-n} + z_{k}$$

where  $x_{k-n} = x_{\ell}(kT - nT)$ .

# Quick summary



• Transmitted signal  $s_{\ell}(t) = \sum_{n=-\infty}^{\infty} I_n g(t - nT)$ 

- Received signal  $r_{\ell}(t) = \sum_{n=-\infty}^{\infty} I_n h_{\ell}(t nT) + n_{\ell}(t)$  with  $h_{\ell}(t) = g(t) \star c_{\ell}(t)$ .
- Matched filter output  $y_{\ell}(t) = \sum_{n=-\infty}^{\infty} I_n x_{\ell}(t nT) + z_{\ell}(t)$  with  $x_{\ell}(t) = h_{\ell}(t) \star h_{\ell}^{\star}(-t)$  or equivalently

$$X_{\ell}(f) = |H_{\ell}(f)|^2 = |G(f)|^2 |C_{\ell}(f)|^2.$$

• Sampling at t = kT and  $y_k = \sum_{n=-\infty}^{\infty} I_n x_{k-n} + z_k$ .

The transmission power is not  $||g(t) \star c_{\ell}(t)||^2$ , which the receiver filter is to match! The additive interference to the transmitted information  $I_k$  is not just  $z_k$  but includes  $\sum_{n=-\infty, n \neq k}^{\infty} I_n x_{k-n}$ . So match filter is good for single transmission but may not be proper for consecutive transmissions!

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# Eye pattern

- A good tool to examine inter-symbol interference is the eye pattern.
  - Eye pattern: The synchronized superposition of all possible realizations of the signal of interest viewed within a particular signaling interval.



Eye pattern for  $r_{\ell}(t)$  for g(t) being half-cycle sine wave with duration  $T_b$ ,  $c_{\ell}(t) = \delta(t)$  and error-free BPSK transmission.



Now with  $c_{\ell}(t) = \delta(t)$ , we have, for example,

$$g(t) = \begin{cases} \frac{1}{\sqrt{T}}, & 0 \le t < T \\ 0, & \text{otherwise} \end{cases}$$
 (Time-limited; hence, band-unlimited!  

$$h_{\ell}(t) = g(t) \star c_{\ell}(t) = g(t)$$

$$x_{\ell}(t) = h_{\ell}(t) \star h_{\ell}^{*}(-t) = \int_{-\infty}^{\infty} h_{\ell}(\tau) h_{\ell}^{*}(-(t-\tau)) d\tau$$

$$= \int_{-\infty}^{\infty} g(\tau) g(\tau-t) d\tau = \begin{cases} \frac{T-|t|}{T}, & |t| \le T \\ 0, & \text{otherwise} \end{cases}$$

From here, you shall know the difference of using  $h_{\ell}^*(T-t)$  and  $h_{\ell}^*(-t)$ , where the former samples at t = T, while the latter samples at t = 0.

## Perfect eye pattern (at Rx)





# Example: BPSK with 1/T = 1K and W = 3K



# Example: BPSK with 1/T = 1K and W = 1K



# Example: BPSK with 1/T = 1K and W = 0.9K



# Example: BPSK with 1/T = 1K and W = 0.8K



# Example: BPSK with 1/T = 1K and W = 0.7K



# Example: BPSK with 1/T = 1K and W = 0.6K



# Example: 4PAM with 1/T = 1K and W = 1K





## Conclusion: The smaller W

- Horizontal: Decision is more sensible to timing error.
- Vertical: Decision is more sensible to noise.

# Example (revisited)

Now changing to  $c_{\ell}(t) = \delta(t) + \delta(t - T)$ , we have, for example,

$$g(t) = \begin{cases} \frac{1}{\sqrt{T}}, & 0 \le t < T\\ 0, & \text{otherwise} \end{cases} \text{ (Time-limited; hence, band-unlimited!)} \\ h_{\ell}(t) = g(t) * c_{\ell}(t) = g(t) + g(t - T) = g(t/2) \\ x_{\ell}(t) = h_{\ell}(t) * \underbrace{h_{\ell}^{*}(2T - t)}_{\text{causal}} = \int_{-\infty}^{\infty} h_{\ell}(\tau)h_{\ell}^{*}(2T - (t - \tau))d\tau \\ = \int_{-\infty}^{\infty} g(\tau/2)g((2T + \tau - t)/2)d\tau \\ = \begin{cases} \frac{2T - |t - 2T|}{T}, & |t - 2T| \le 2T \\ 0, & \text{otherwise} \end{cases}$$

$$y_{k} = y_{\ell}(kT) = \sum_{n=-\infty} I_{n} x_{\ell}(kT - nT) + z_{\ell}(kT) = 2I_{k-2} + \underbrace{I_{k-1} + I_{k-3}}_{|S|} + z_{k}$$

The question next to be asked is that how to remove the ISI?

9.2-1 Design of band-limited signals for no intersymbol interference -The Nyquist criterion

## Intersymbol interference channel

$$y_k = \sum_{n=-\infty}^{\infty} I_n x_{k-n} + z_k$$

### Design goal

Given W, T and  $c_{\ell}(t)$ , we would like to design g(t) such that

$$x_{k-n} = \delta_{k-n}$$

and that

$$y_k = \sum_{n=-\infty}^{\infty} I_n x_{k-n} + z_k = I_k + z_k$$

Here,  $\delta_k$  is the Kronecker delta function, defined as

$$\delta_k = \begin{cases} 1 & k = 0 \\ 0 & k \neq 0 \end{cases}$$

### Theorem 1 (Nyquist criterion)

Let  $x_{\ell}(t)$  be a band-limited signal with band [-W, W] and  $x_{\ell}(0) = 1$ . Sample at rate  $\frac{1}{T}$  such that

$$x_k = x_\ell(kT) = \int_{-\infty}^{\infty} x_\ell(t)\delta(t-kT) dt = \delta_k$$

if and only if

$$\sum_{m=-\infty}^{\infty} X_{\ell}\left(f-\frac{m}{T}\right) = T$$

where  $X_{\ell}(f) = \mathcal{F}\{x_{\ell}(t)\}.$ 

#### **Proof:**

The condition of  $x_k = x_\ell(kT) = \delta_k$  is equivalent to

$$x_{\ell}(t)\sum_{k=-\infty}^{\infty}\delta(t-kT) = \sum_{k=-\infty}^{\infty}x_{\ell}(kT)\delta(t-kT) = \delta(t).$$
(1)

We however have (from Lemma 2 in the next slide)

$$\mathcal{F}\left\{\sum_{k=-\infty}^{\infty}\delta(t-kT)\right\} = \frac{1}{T}\sum_{m=-\infty}^{\infty}\delta\left(f-\frac{m}{T}\right) \text{ and } \mathcal{F}\left\{\delta(t)\right\} = 1.$$

Taking Fourier transform on both sides of (1) gives

$$X_{\ell}(f) \star \left[\frac{1}{T}\sum_{m=-\infty}^{\infty}\delta\left(f-\frac{m}{T}\right)\right] = \frac{1}{T}\sum_{m=-\infty}^{\infty}X_{\ell}\left(f-\frac{m}{T}\right) = 1,$$

and proves the theorem.

# Proof of key lemma

## Lemma 2

$$\mathcal{F}\left\{\sum_{k=-\infty}^{\infty}\delta(t-kT)\right\} = \frac{1}{T}\sum_{m=-\infty}^{\infty}\delta\left(f-\frac{m}{T}\right)$$

**Proof:** Consider the function

$$\alpha(f) = \frac{1}{T} \sum_{m=-\infty}^{\infty} \delta\left(f - \frac{m}{T}\right)$$

which is periodic with period  $\frac{1}{T}$ . From Fourier series, we have

$$\alpha(f) = \frac{1}{1/T} \sum_{n=-\infty}^{\infty} c_n e^{-i 2\pi n f / (1/T)} = \sum_{n=-\infty}^{\infty} e^{-i 2\pi n f T}$$

where

$$c_n = \int_{-\frac{1}{2T}}^{\frac{1}{2T}} \alpha(f) e^{i 2\pi n f/(1/T)} df = \int_{-\frac{1}{2T}}^{\frac{1}{2T}} \frac{1}{T} \delta(f) e^{i 2\pi n f/(1/T)} df = \frac{1}{T}$$

by following the replication property of  $\delta(f)$ .

Next by linearity of Fourier transform, we have

$$\mathcal{F}\left\{\sum_{k=-\infty}^{\infty}\delta(t-kT)\right\} = \int_{-\infty}^{\infty}\sum_{k=-\infty}^{\infty}\delta(t-kT)e^{-i2\pi ft} dt$$
$$= \sum_{k=-\infty}^{\infty}e^{-i2\pi fkT}.$$

# Implication of Nyquist criterion

$$x_k = \delta_k \quad \text{iff} \quad \sum_{m=-\infty}^{\infty} X_\ell \left( f - \frac{m}{T} \right) = T$$
 (2)

• 
$$2W < \frac{1}{T}$$
:  
 $X_{\ell} \left( f - \frac{m}{T} \right)$  and  $X_{\ell} \left( f - \frac{m'}{T} \right)$  do not overlap for  $m \neq m'$ .  
Hence (2) is impossible!

**2** 
$$W = \frac{1}{T}$$
:  
This means  $X_{\ell}(f) = T \operatorname{rect}(Tf)$ , hence  $x_{\ell}(t) = \operatorname{sinc}(\frac{t}{T})$ .  
Theoretically ok but physically impossible!

3 So we need 
$$2W > \frac{1}{T}$$

## Channel bandwidth must be larger than sampling rate.

$$x_k = \delta_k$$
 iff  $\sum_{m=-\infty}^{\infty} X_\ell \left( f - \frac{m}{T} \right) = T$ 



## Definition 1

$$X_{rc}(f) = \begin{cases} T, & 0 \le |f| \le \frac{1-\beta}{2T} \\ \frac{T}{2} \left\{ 1 + \cos\left[\frac{\pi T}{\beta} \left(|f| - \frac{1-\beta}{2T}\right)\right] \right\}, & \frac{1-\beta}{2T} \le |f| \le \frac{1+\beta}{2T} \\ 0, & otherwise \end{cases}$$

where  $0 \le \beta \le 1$  is the roll-off factor.

$$x_{rc}(t) = sinc\left(\frac{t}{T}\right) \frac{\cos\left(\pi\beta t/T\right)}{1 - 4\beta^2 t^2/T^2}$$



$$X_{rc}\left(f-\frac{m}{T}\right)$$
 with  $m=-1,0,1$ 









## Assuming

$$C_{\ell}(f) = 1$$

• Let 
$$G_T(f) = \mathcal{F}\{g(t)\}.$$

• Let 
$$G_R(f) = \mathcal{F}\{h_{\ell}^*(-t)\}$$
, where  $h_{\ell}(t) = g(t) \star c_{\ell}(t) = g(t)$ .

This gives

$$G_R(f) = \mathcal{F}\{g^*(-t)\} = G_T^*(f)$$

 $\mathsf{and}$ 

$$X_{rc}(f) = G_T(f)C_\ell(f)G_R(f) = G_T(f)G_R(f) = |G_T(f)|^2.$$

• Note  $|x_{rc}(t)| \approx \frac{1}{t^3}$  for large |t|. So it can be truncated at  $\pm t_0$  for some  $t_0$  large.

$$x_{rc}(t) = \operatorname{sinc}\left(\frac{t}{T}\right) \frac{\cos\left(\pi\beta t/T\right)}{1-4\beta^2 t^2/T^2}$$

So we can set

$$G_T(f) = \sqrt{X_{rc}(f)}e^{-\imath 2\pi f t_0}$$
 and  $G_R(f) = G_T^*(f)$ .

# 9.2-2 Design of band-limited signals with controlled ISI: Partial response signals

# Duobinary pulse

We relax the no-ISI condition so that

$$x_k = \delta_k + \delta_{k-1} \implies y_k = I_k + I_{k-1} + z_k$$

Following similar arguments (see the next slide), it shows

$$\frac{1}{T}\sum_{m=-\infty}^{\infty}X_{\ell}\left(f-\frac{m}{T}\right)=1+e^{-\imath 2\pi fT}$$

Setting  $2W = \frac{1}{T}$  (with W being the bandwidth of  $X_{\ell}(f)$ ), we have

$$X_{\ell}(f) = T\left(1 + e^{-i2\pi fT}\right) \operatorname{rect}\left(\frac{f}{2W}\right)$$
$$x_{\ell}(t) = \operatorname{sinc}\left(2Wt\right) + \operatorname{sinc}\left[2\left(Wt - \frac{1}{2}\right)\right]$$

## This is called duobinary pulse.

**Proof:** The condition of  $x_k = \delta_k + \delta_{k-1}$  is equivalent to

$$x_{\ell}(t)\sum_{k=-\infty}^{\infty}\delta(t-kT) = \sum_{k=-\infty}^{\infty}x_{\ell}(kT)\delta(t-kT) = \delta(t) + \delta(t-T).$$
(3)

We however have

$$\mathcal{F}\left\{\sum_{k=-\infty}^{\infty}\delta(t-kT)\right\} = \frac{1}{T}\sum_{m=-\infty}^{\infty}\delta\left(f-\frac{m}{T}\right) \text{ and } \mathcal{F}\left\{\delta(t)\right\} = 1.$$

Taking Fourier transform on both sides of (3) gives

$$X_{\ell}(f) \star \left[\frac{1}{T} \sum_{m=-\infty}^{\infty} \delta\left(f - \frac{m}{T}\right)\right] = \frac{1}{T} \sum_{m=-\infty}^{\infty} X_{\ell}\left(f - \frac{m}{T}\right) = 1 + e^{-i2\pi fT},$$

and proves the theorem.

When 
$$W = \frac{1}{2T}$$
,  
 $X_{\ell}(f) = \begin{cases} T(1 + e^{-i2\pi fT}) \operatorname{rect}\left(\frac{f}{2W}\right), & \text{with controlled ISI}\\ T \operatorname{rect}\left(\frac{f}{2W}\right), & \text{without controlled ISI} \end{cases}$ 

$$X_{\ell}(f) \text{ with controlled SI} \xrightarrow{|X(f)|}_{W} \xrightarrow{\frac{1}{W} \cos \frac{\pi f}{2W}}_{-W} = K$$

The duobinary filter is apparently more physically realizable!

Recall that with no controlled ISI, we require  $W > \frac{1}{2T}$  because channel bandwidth  $W = \frac{1}{2T}$  is not physically realizable. But now with controlled ISI,  $W = \frac{1}{2T}$  becomes physically realizable.

# 9.2-3 Data detection for controlled ISI

# Precoding for duobinary pulses

Received signal for duobinary shaping is

$$y_k = I_k + I_{k-1} + z_k.$$

We could **precode** the sequence  $\{I_k\}$  to simplify detection.

Example 1 (Binary case - Differential encoding) Given the binary bit (information) stream  $\{b_k\}$ • Define  $P_k = b_k \oplus P_{k-1} \in \{0, 1\}$ .

• Set 
$$I_k = 2P_k - 1 \in \{\pm 1\}.$$

• Receive 
$$y_k = I_k + I_{k-1} + z_k$$
.

$$I_k + I_{k-1} = \begin{cases} \pm 2, & \text{if } b_k = 0\\ 0, & \text{if } b_k = 1. \end{cases}$$

•  $\hat{b}_k = 0 \ if |y_k| > 1 \ and \ \hat{b}_k = 1 \ if |y_k| \le 1$ 

# Summary of precoding system

Compare 
$$\begin{cases} \frac{1}{T} \sum_{m=-\infty}^{\infty} X_{\ell} \left( f - \frac{m}{T} \right) = 1\\ \frac{1}{T} \sum_{m=-\infty}^{\infty} X_{\ell}^{(p)} \left( f - \frac{m}{T} \right) = 1 + e^{-i 2\pi f T}. \end{cases}$$

We can say  $X_{\ell}^{(p)}(f) = X_{\ell}(f) \left(1 + e^{-\imath 2\pi fT}\right)$  and

$$\begin{cases} y_{\ell}(t) = \sum_{n=-\infty}^{\infty} I_n x_{\ell}(t - nT) + z_{\ell}(t) \\ y_{\ell}^{(p)}(t) = \sum_{n=-\infty}^{\infty} I_n x_{\ell}^{(p)}(t - nT) + z_{\ell}(t) \end{cases}$$

and

$$\begin{cases} y_k = I_k + z_k \\ y_k^{(p)} = I_k + I_{k-1} + z_k \end{cases} \Rightarrow \begin{cases} y_k \leq 0 \\ |y_k^{(p)}| \leq 1 \end{cases}$$

Note that  $||x_{\ell}(t)||^2$  does not decide the transmission energy!

# Summary of precoding system

$$\Rightarrow \begin{cases} P_e = Q\left(\sqrt{\frac{1}{\mathbb{E}[z_k^2]}}\right) = Q\left(\sqrt{\frac{\mathcal{E}_{b,\ell}/\|g_T(t)\|^2}{N_0\|g_R(t)\|^2}}\right), & \text{error rate for } I_k \\ = Q\left(\sqrt{\frac{2\mathcal{E}_b}{N_0}}\frac{1}{\|g_T(t)\|^2\|g_R(t)\|^2}\right) = Q(\sqrt{2a\gamma_b}) \\ P_e^{(p)} = \frac{3}{2}Q(\sqrt{2a\gamma_b}) - \frac{1}{2}Q(3\sqrt{2a\gamma_b}), & \text{error rate for } b_k \end{cases}$$

where  

$$\begin{cases}
I_k \in \{\pm 1\} \\
\text{the (lowpass) transmission energy per bit } \mathcal{E}_{b,\ell} = \|g_T(t)\|^2 = 2\mathcal{E}_b \\
\text{the lowpass noise } \mathbb{E}[z_k^2] = \sigma_\ell^2 \|g_R(t)\|^2 \\
\text{with } n_\ell(t) \text{ having two-sided PSD } \sigma_\ell^2 = N_0 \\
a = \frac{1}{\|g_T(t)\|^2 \|g_R(t)\|^2} \text{ (subject to } g_T(t) \star c_\ell(t) \star g_R(t) = x_\ell(t))
\end{cases}$$

Hence, pre-coding technique provides a better spectrum efficiency at the price of performance degradation.

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# 9.2-4 Signal design for channel with distortion

In general, we have  $C_{\ell}(f) \neq \operatorname{rect}\left(\frac{f}{2W}\right)$ , and in this case

$$s_{\ell}(t) = \sum_{n=-\infty}^{\infty} I_n g_T(t - nT)$$
  

$$r_{\ell}(t) = c_{\ell}(t) \star s_{\ell}(t) + n_{\ell}(t)$$
  

$$y_{\ell}(t) = r_{\ell}(t) \star g_R(t) = s_{\ell}(t) \star c_{\ell}(t) \star g_R(t) + z_{\ell}(t)$$
  

$$y_k = y_{\ell}(kT) = I_k + z_k \quad (We hope there is no ISI!)$$

where

- $g_T(t)$  transmit filter
- $c_{\ell}(t)$  channel impulse response
- $g_R(t)$  receive filter

### Hence

$$\begin{cases} x_{rc}(t) = g_{T}(t) \star c_{\ell}(t) \star g_{R}(t) \\ X_{rc}(f) = G_{T}(f)C_{\ell}(f)G_{R}(f) \\ S_{z_{\ell}}(f) = N_{0}|G_{R}(f)|^{2} \text{ because we assume } \{I_{k}\} \text{ real} \end{cases}$$

If  $c_{\ell}(t)$  is known to Tx, then we may choose to "pre-equalize" the channel effect at Tx:

$$|G_T(f)| = \frac{\sqrt{X_{rc}(f)}}{|C_\ell(f)|}$$
 and  $|G_R(f)| = \sqrt{X_{rc}(f)}$ 

Then, ISI is avoided; also, the noise power remains

$$\mathbb{E}[z_k^2] = \int_{-\infty}^{\infty} \underbrace{S_{z_\ell}(f)}_{=N_0|G_R(f)|^2} df = N_0 \int_{-\infty}^{\infty} X_{rc}(f) df = N_0 x_{rc}(0) = N_0.$$

## $I_n \in \{\pm d\}$ and $X_{rc}(f) = G_T(f)C_\ell(f)G_R(f)$

Signal power

$$P_{av,\ell} = \frac{d^2 \|g_T(t)\|^2}{T} = \frac{d^2}{T} \int_{-\infty}^{\infty} \frac{X_{rc}(f)}{|C_{\ell}(f)|^2} df$$

Error probability

$$P_{b} = Q\left(\sqrt{\frac{d_{12}^{2}}{4\mathbb{E}[z_{k}^{2}]}}\right)$$
$$= Q\left(\frac{d}{\sqrt{\mathbb{E}[z_{k}^{2}]}}\right) = Q\left(\sqrt{\frac{P_{av,\ell}T}{N_{0}\left[\int_{-\infty}^{\infty}\frac{X_{rc}(f)}{|C_{\ell}(f)|^{2}}df\right]}}\right)$$

# If $c_{\ell}(t)$ only known to Rx

We can only equalize the "channel effect" at Rx:

$$|G_T(f)| = \sqrt{X_{rc}(f)}$$
 and  $|G_R(f)| = \frac{\sqrt{X_{rc}(f)}}{|C_\ell(f)|}$ .

Signal power

$$P_{av,\ell} = \frac{d^2}{T} \|g_T(t)\|^2 = \frac{d^2}{T} \int_{-\infty}^{\infty} X_{rc}(f) df = \frac{d^2}{T}$$

Noise power

$$\mathbb{E}[z_k^2] = \int_{-\infty}^{\infty} S_{z_\ell}(f) df = N_0 \int_{-\infty}^{\infty} \frac{X_{rc}(f)}{|C_\ell(f)|^2} df$$

Error probability

$$P_b = Q\left(\frac{d}{\sqrt{\mathbb{E}[z_k^2]}}\right) = Q\left(\sqrt{\frac{P_{av,\ell}T}{N_0\left[\int_{-\infty}^{\infty}\frac{X_{rc}(f)}{|C_\ell(f)|^2}df\right]}}\right)$$

# If however $c_{\ell}(t)$ known to both Tx and Rx

We may design:

$$|G_T(f)| = \sqrt{\frac{X_{rc}(f)}{|C_\ell(f)|}}$$
 and  $|G_R(f)| = \sqrt{\frac{X_{rc}(f)}{|C_\ell(f)|}}$ 

Signal power

$$P_{av,\ell} = \frac{d^2}{T} \|g_T(t)\|^2 = \frac{d^2}{T} \int_{-\infty}^{\infty} \frac{X_{rc}(f)}{|C_\ell(f)|} df$$

Noise power

$$\mathbb{E}[z_k^2] = \int_{-\infty}^{\infty} S_{z_\ell}(f) df = N_0 \int_{-\infty}^{\infty} \frac{X_{rc}(f)}{|C_\ell(f)|} df$$

Error probability

$$P_b = Q\left(\frac{d}{\sqrt{\mathbb{E}[z_k^2]}}\right) = Q\left(\sqrt{\frac{P_{av,\ell}T}{N_0\left[\int_{-\infty}^{\infty}\frac{X_{rc}(f)}{|C_\ell(f)|}df\right]^2}}\right)$$

Either Tx or Rx knows  $c_{\ell}(t)$ 

$$P_{b,T} = P_{b,R} = Q\left(\sqrt{\frac{P_{av,\ell}T}{N_0 \left[\int_{-\infty}^{\infty} \frac{X_{rc}(f)}{|C_\ell(f)|^2} df\right]}}\right)$$

Both Tx and Rx know  $c_{\ell}(t)$ 

$$P_{b,TR} = Q\left(\sqrt{\frac{P_{av,\ell}T}{N_0 \left[\int_{-\infty}^{\infty} \frac{X_{rc}(f)}{|C_\ell(f)|} df\right]^2}}\right)$$

Note from Cauchy-Schwartz inequality

$$\left[\int_{-\infty}^{\infty} \frac{X_{rc}(f)}{|C_{\ell}(f)|} df\right]^{2} = \left\| \left( \sqrt{X_{rc}(f)}, \frac{\sqrt{X_{rc}(f)}}{|C_{\ell}(f)|} \right) \right\|^{2}$$
$$\leq \left\| \sqrt{X_{rc}(f)} \right\|^{2} \left\| \frac{\sqrt{X_{rc}(f)}}{|C_{\ell}(f)|} \right\|^{2} = \int_{-\infty}^{\infty} \frac{X_{rc}(f)}{|C_{\ell}(f)|^{2}} df$$

This shows  $P_{b,TR} \leq P_{b,T} = P_{b,R}$ . "=" holds iff  $|C_{\ell}(f)| = 1$ .

# What you learn from Chapter 9



- Match filter to input pulse shaping and channel impulse response
  - (Good to know) Eye pattern to examine ISI
- Nyquist criterion
  - Sampling rate < channel bandwidth for no ISI (I.e., increasing sampling rate will give more samples for perhaps better performance, but adjacent samples will be eventually "interfered" to each other)
  - Since ISI is unavoidable for high sampling rate, let's accept and face it, and just use controlled ISI.
- A better performance is resulted when both Tx and Rx know the channel.